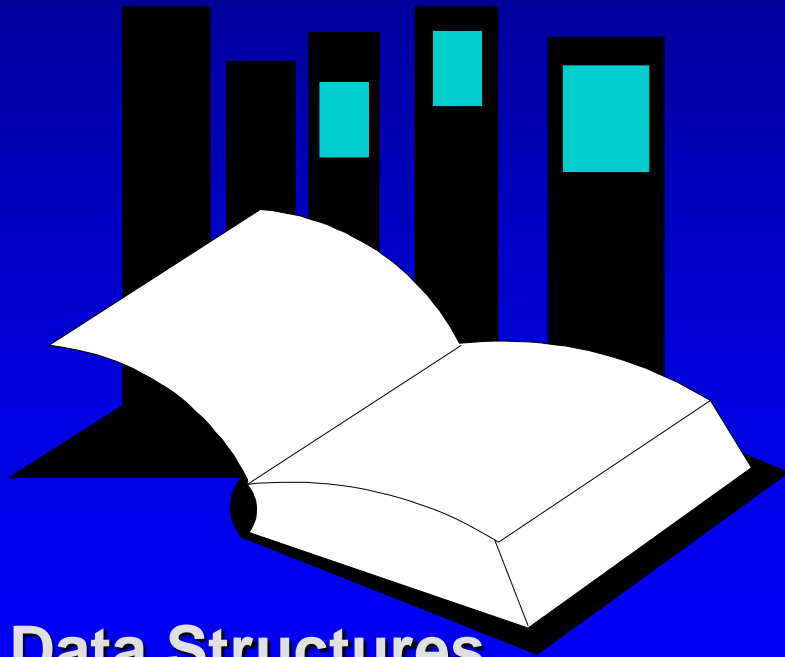
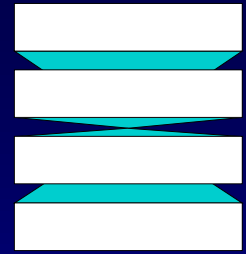


# Using a Stack

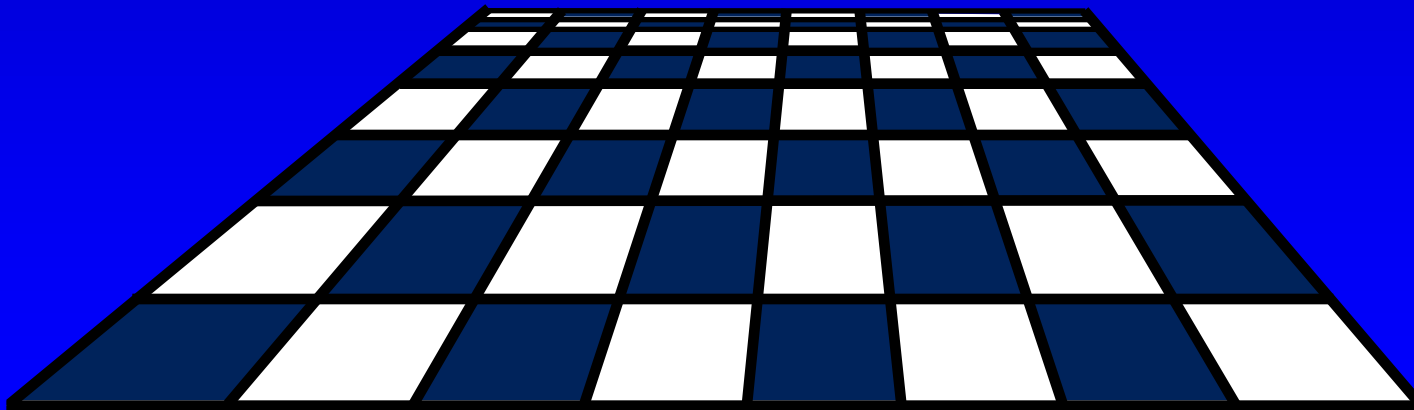
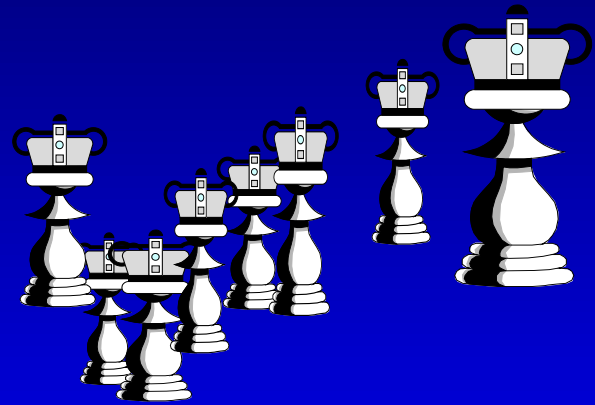


Data Structures  
and Other Objects  
Using Java

- ❑ Chapter 6 introduces the stack data type.
- ❑ Several example applications of stacks are given in that chapter.
- ❑ This presentation shows another use called backtracking to solve the N-Queens problem.

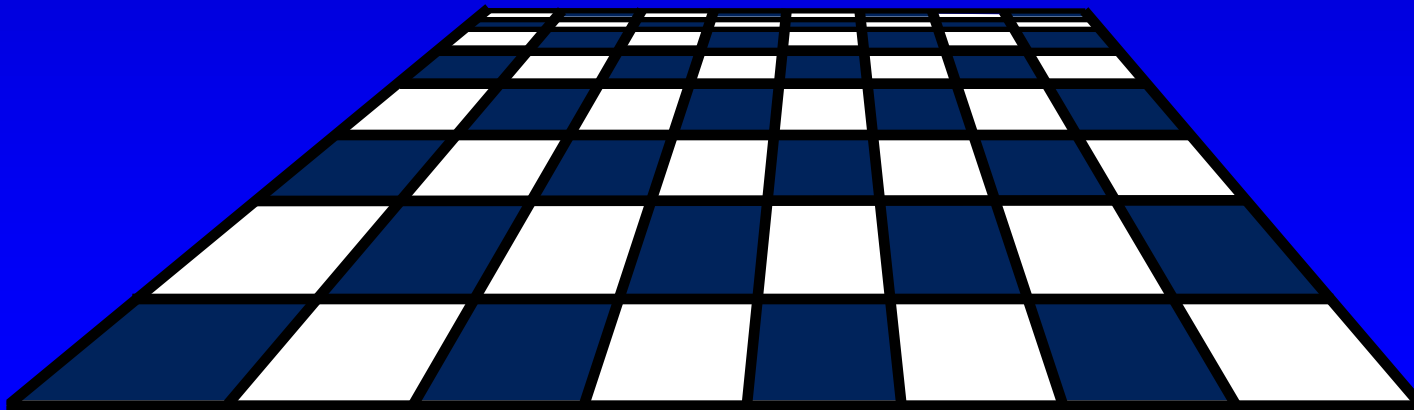
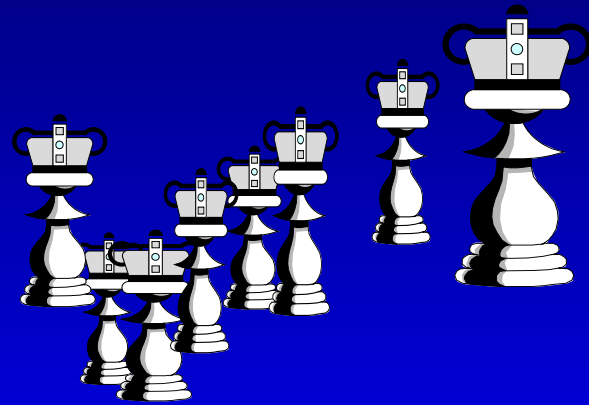
# The N-Queens Problem

- Suppose you have 8 chess queens...
- ...and a chess board



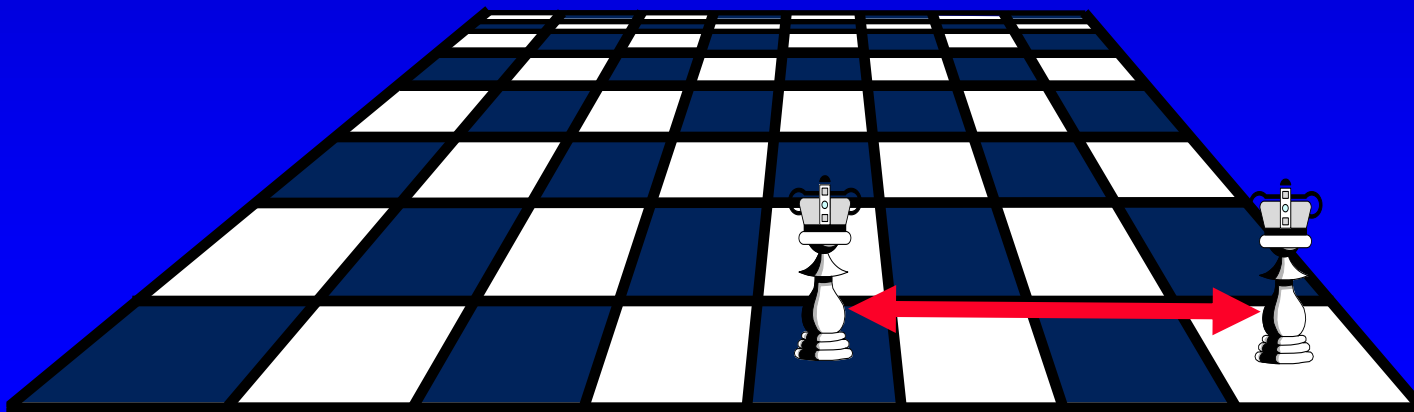
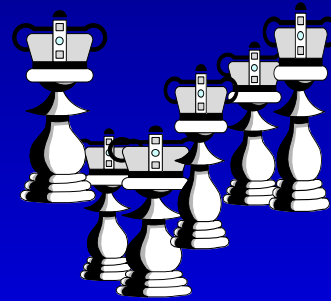
# The N-Queens Problem

*Can the queens be placed on the board so that no two queens are attacking each other ?*



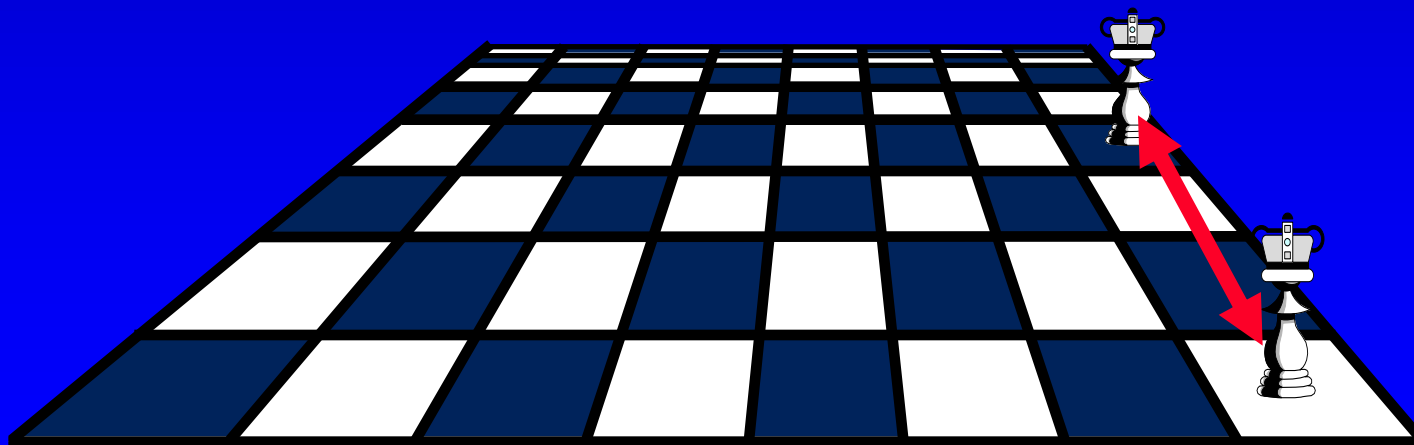
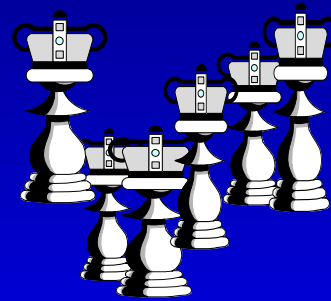
# The N-Queens Problem

Two queens are not  
allowed in the same  
row...



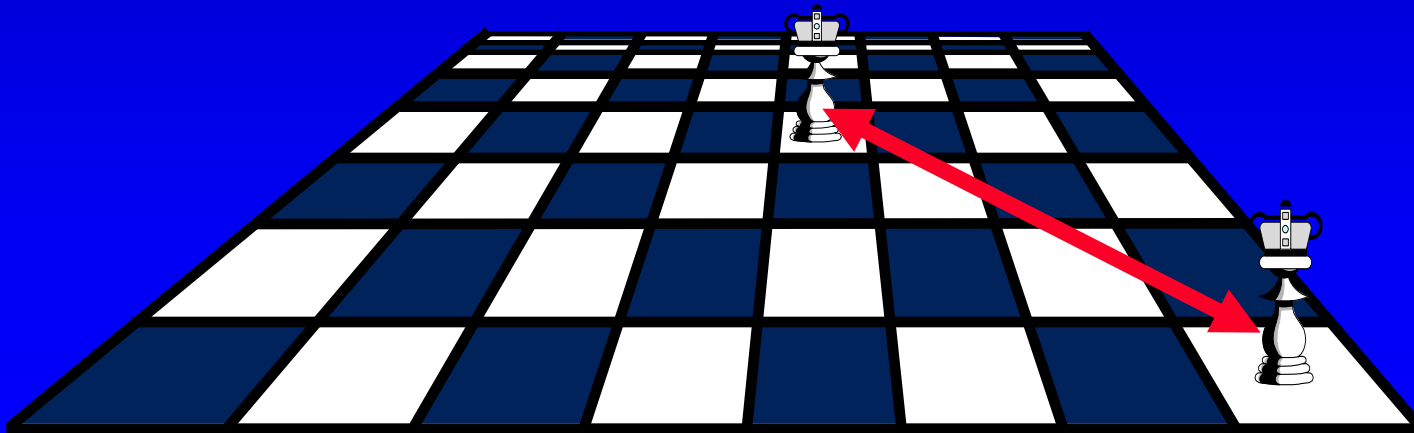
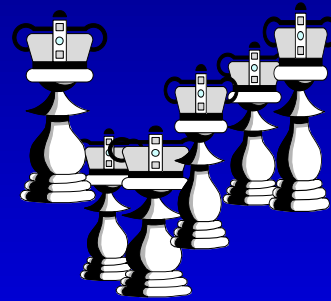
# The N-Queens Problem

Two queens are not allowed in the same row, or in the same column...



# The N-Queens Problem

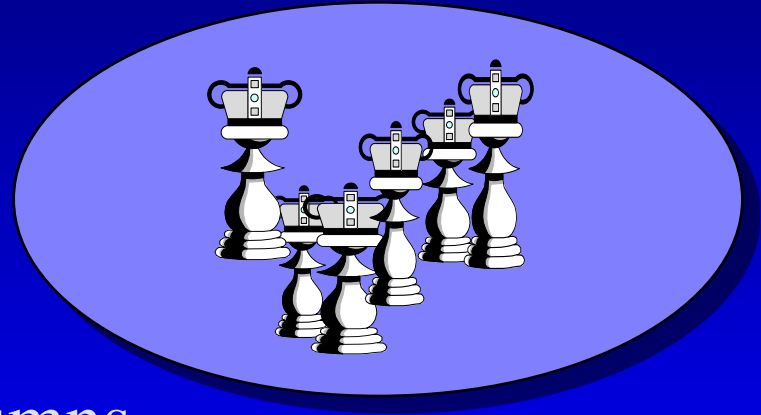
Two queens are not allowed in the same row, or in the same column, or along the same diagonal.



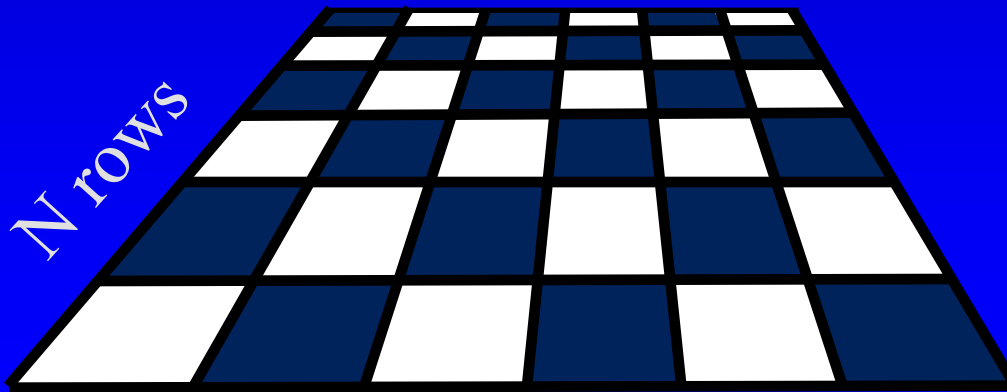
# The N-Queens Problem

The number of queens,  
and the size of the board  
can vary.

N Queens



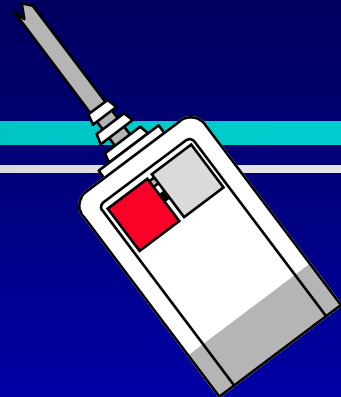
N columns



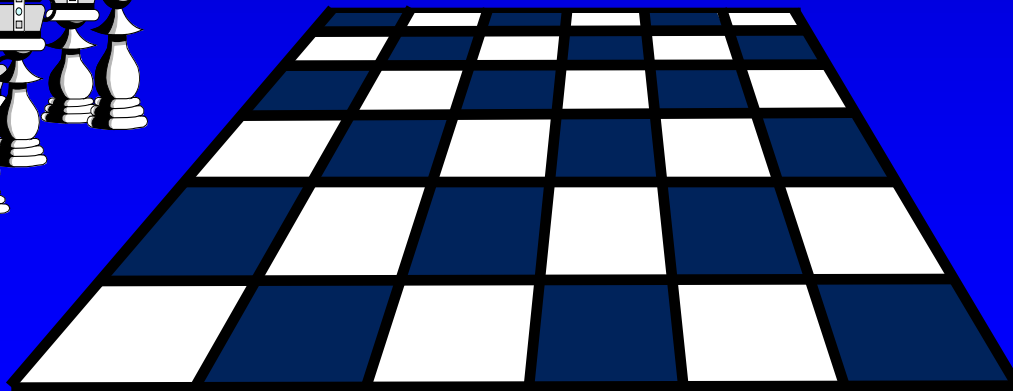
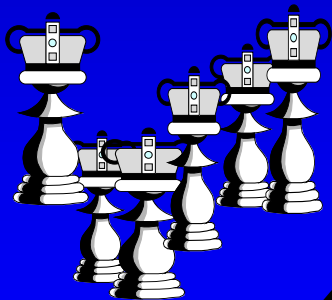
N rows

# The N-Queens Problem

We will write a program which tries to find a way to place  $N$  queens on an  $N \times N$  chess board.



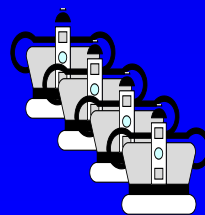
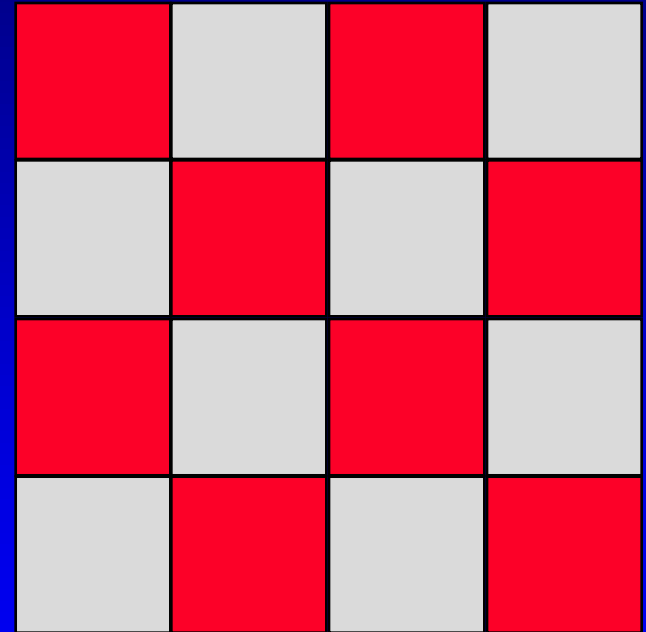
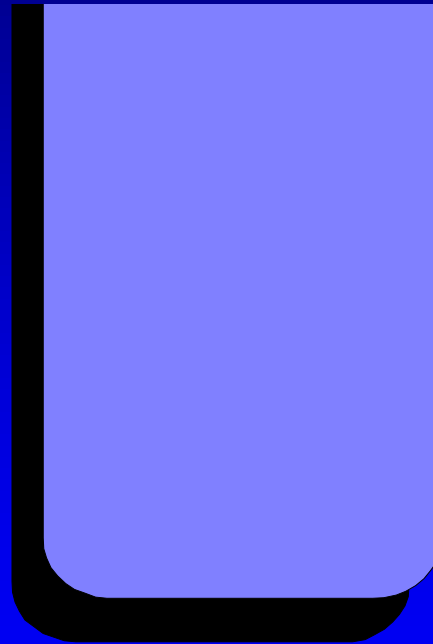
If you can run ega or vga graphics, you can double click on this icon with the left mouse button:





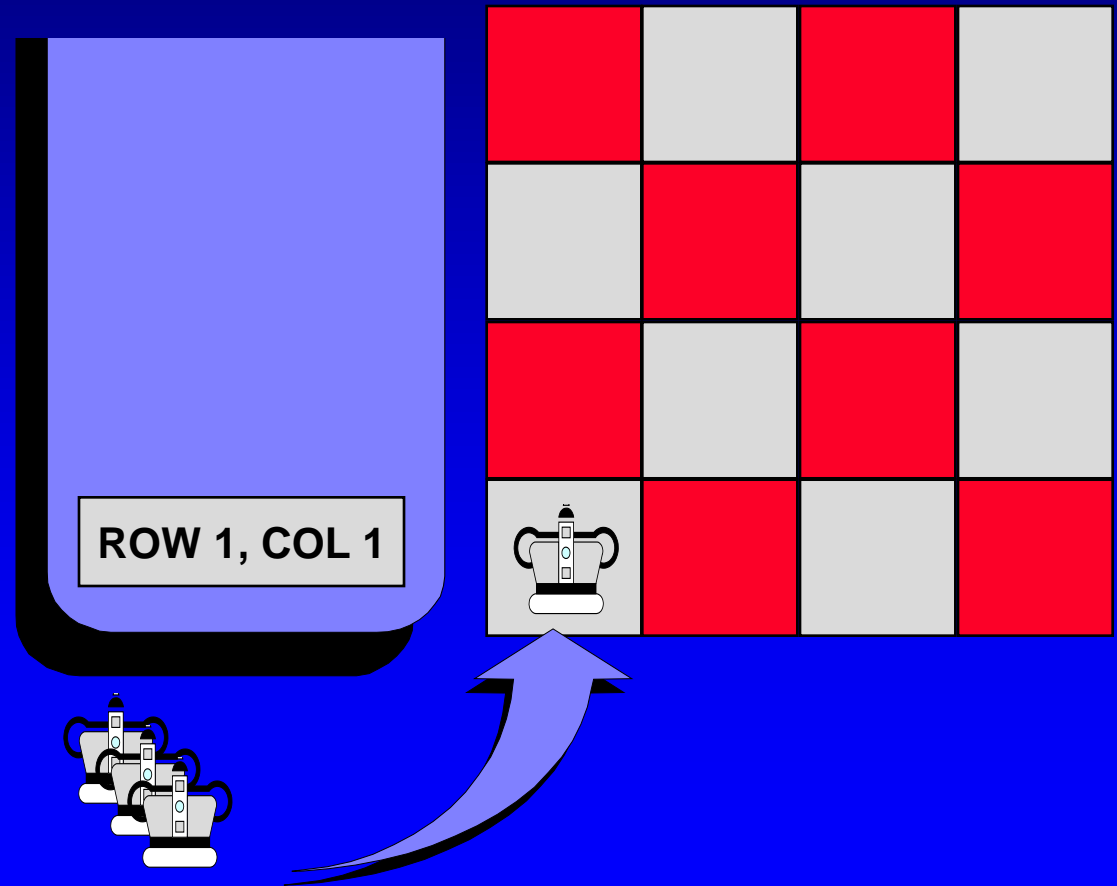
# How the program works

The program  
uses a stack to  
keep track of  
where each  
queen is placed.



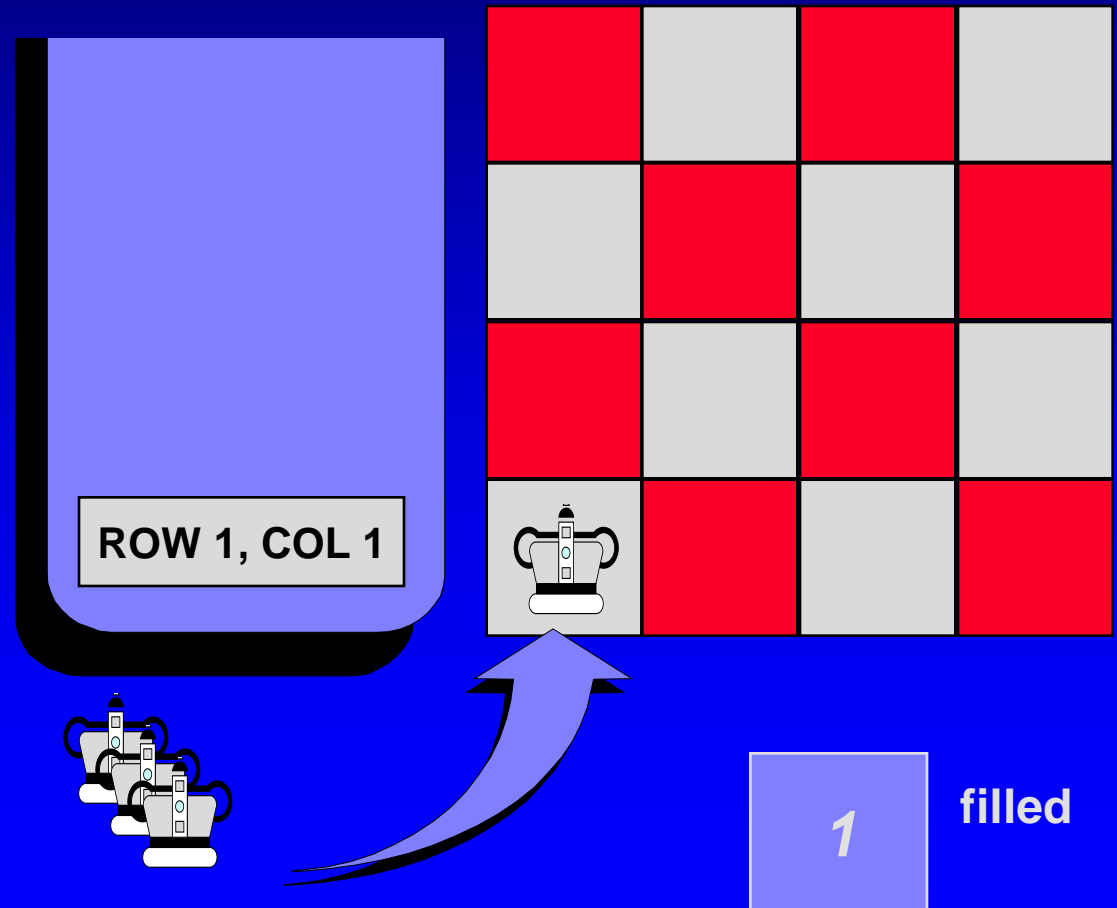
# How the program works

Each time the program decides to place a queen on the board, the position of the new queen is stored in a record which is placed in the stack.



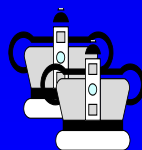
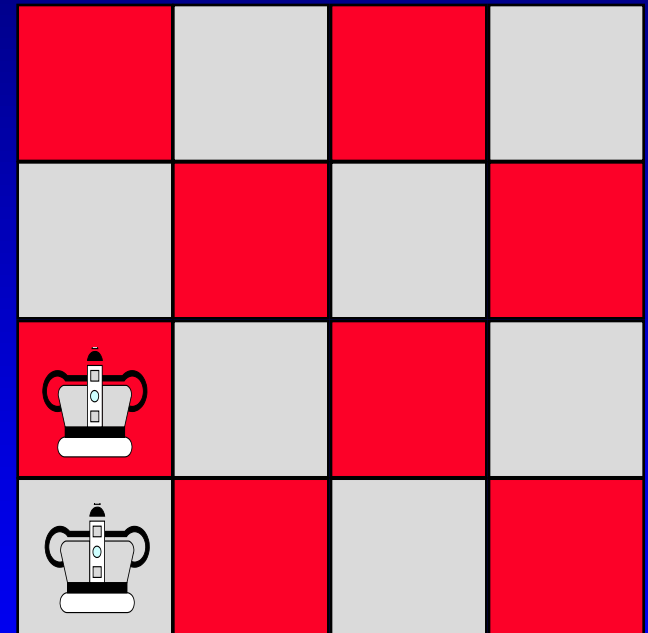
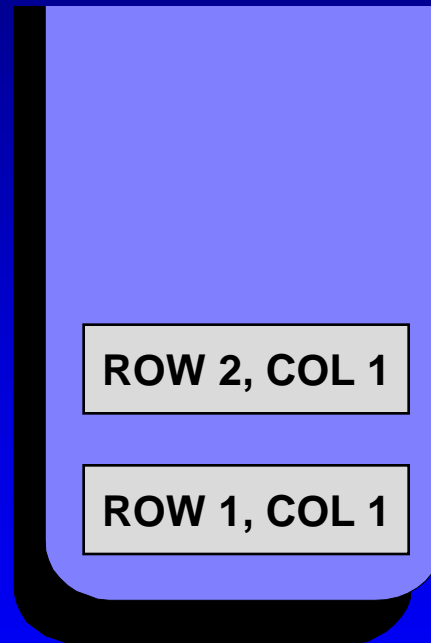
# How the program works

We also have an integer variable to keep track of how many rows have been filled so far.



# How the program works

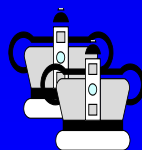
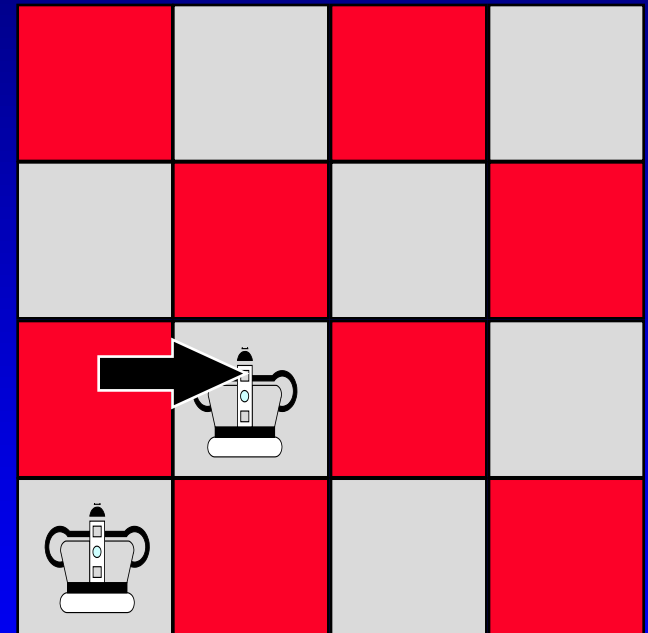
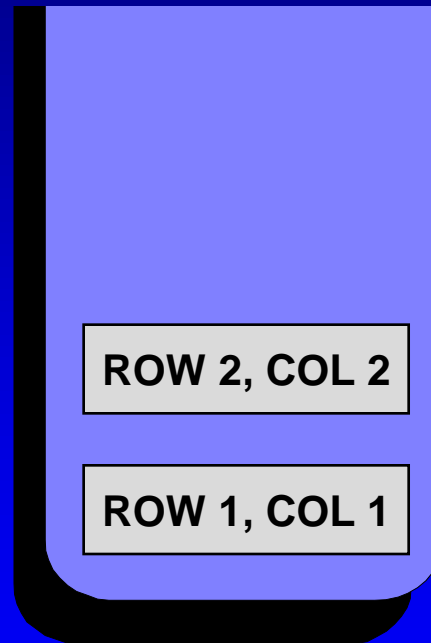
Each time we try to place a new queen in the next row, we start by placing the queen in the first column...



filled

# How the program works

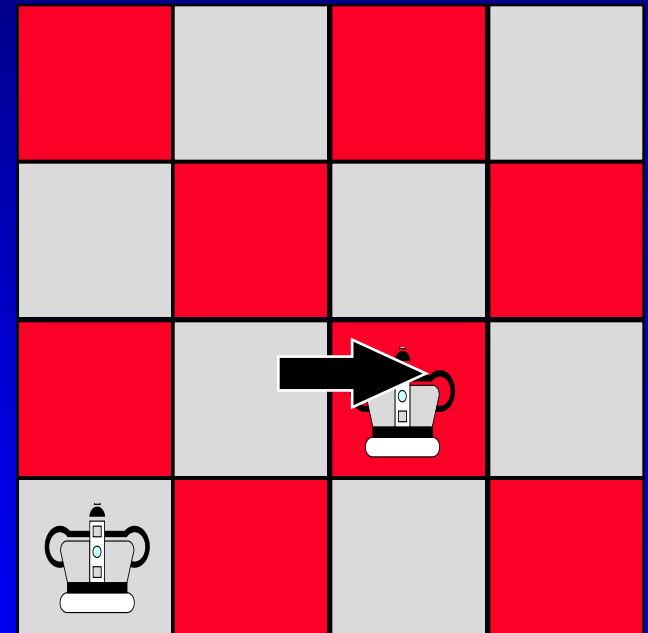
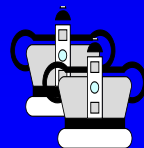
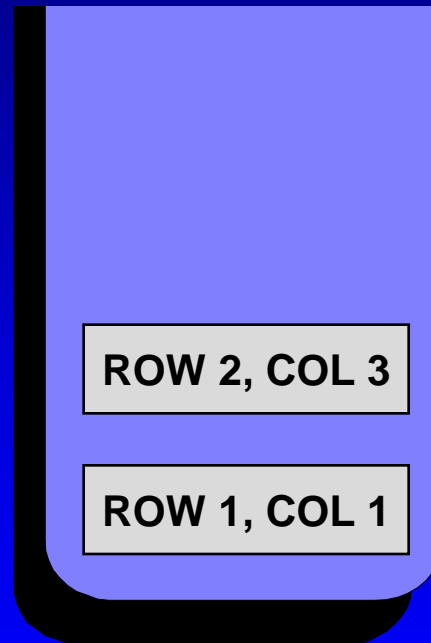
...if there is a conflict with another queen, then we shift the new queen to the next column.



filled

# How the program works

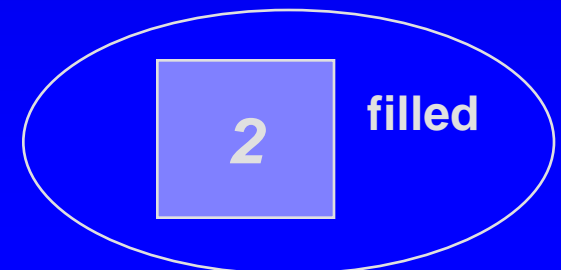
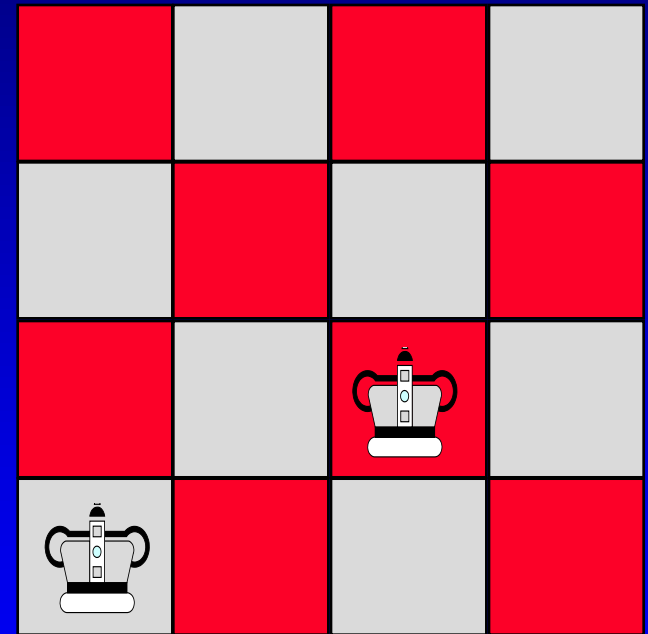
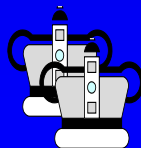
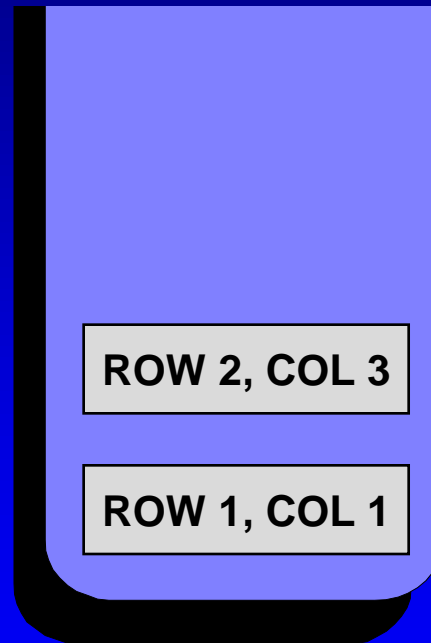
If another conflict occurs, the queen is shifted rightward again.



filled

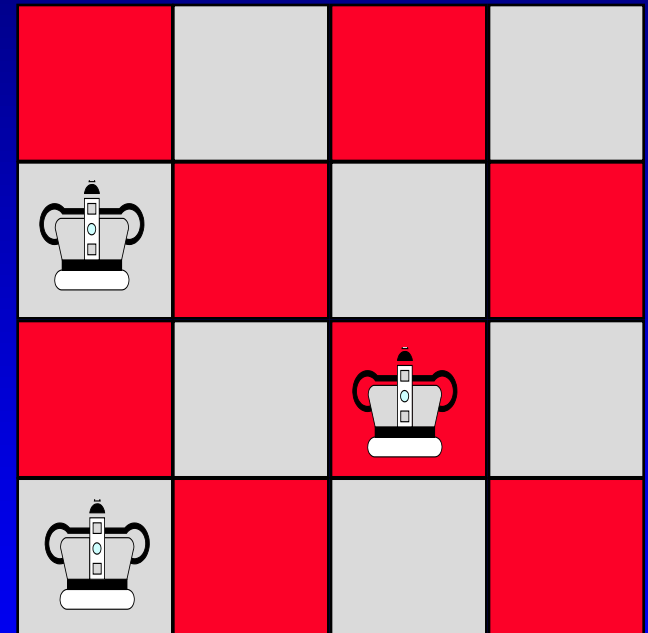
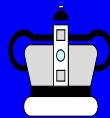
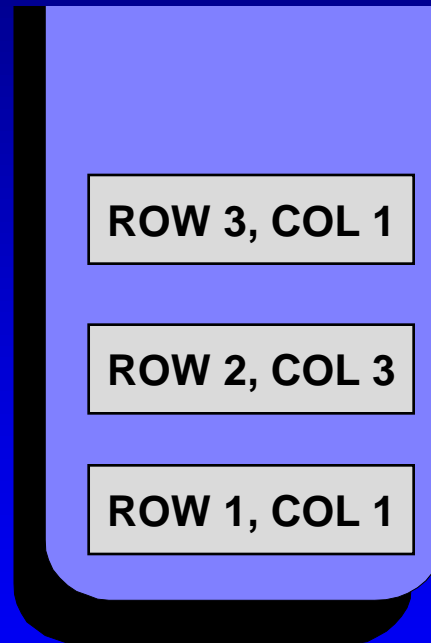
# How the program works

When there are no conflicts, we stop and add one to the value of filled.



# How the program works

Let's look at the third row. The first position we try has a conflict...

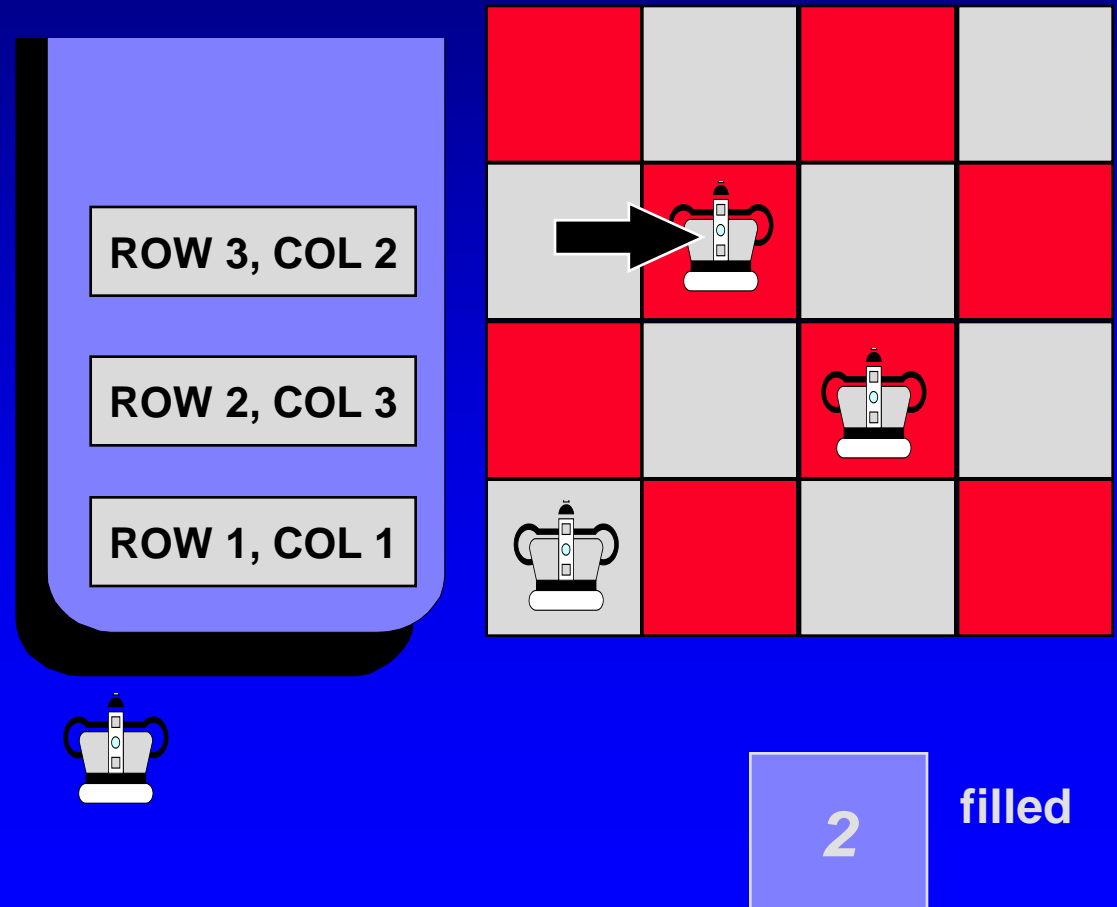


filled



# How the program works

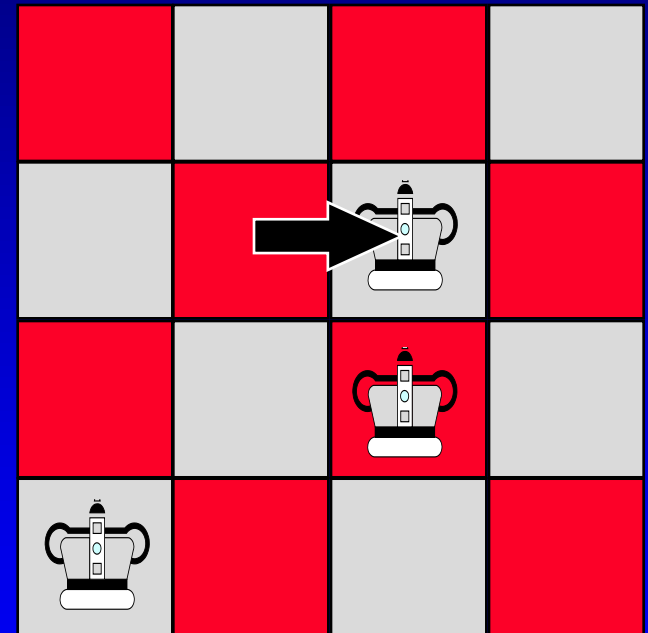
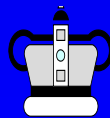
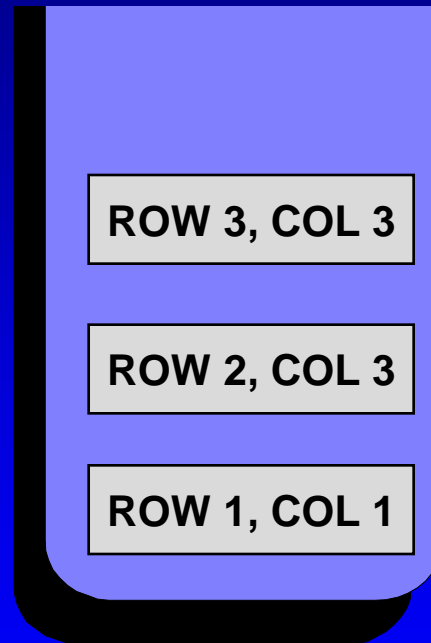
...so we shift to column 2. But another conflict arises...



# How the program works

...and we shift to  
the third column.

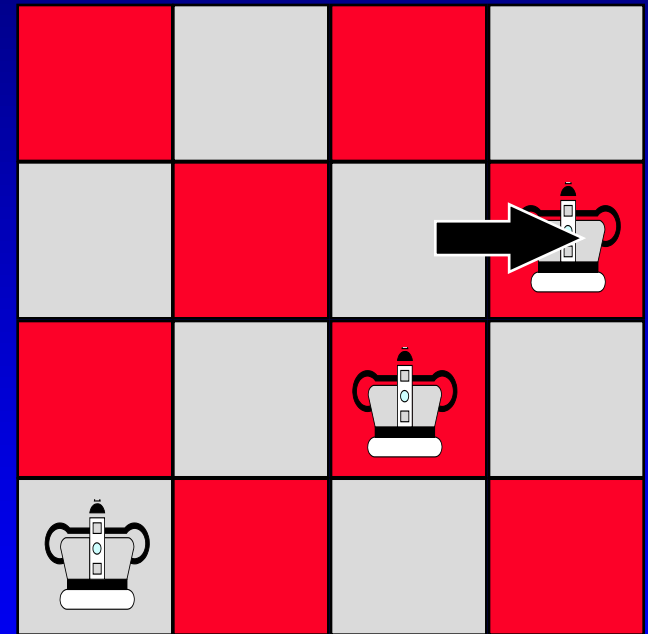
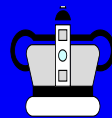
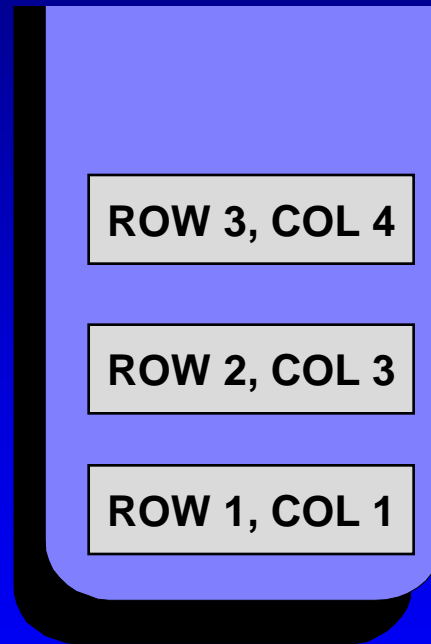
Yet another  
conflict arises...



filled

# How the program works

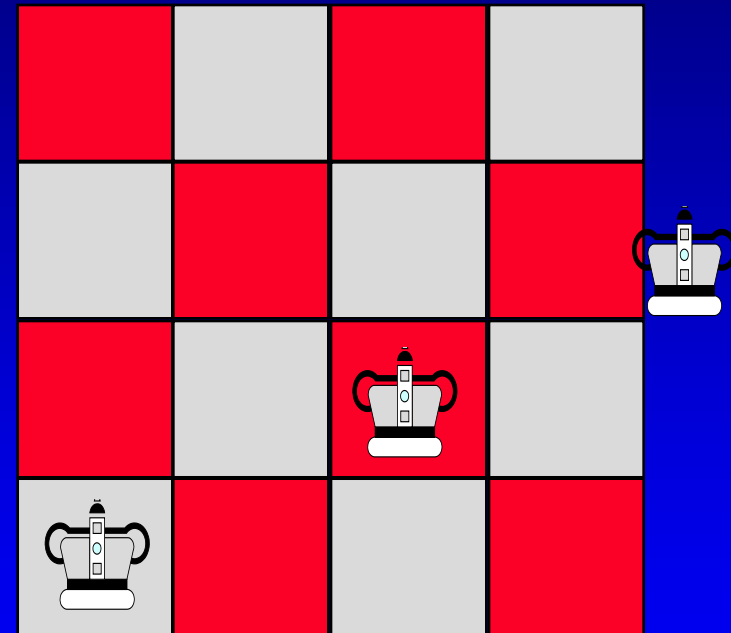
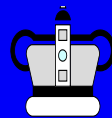
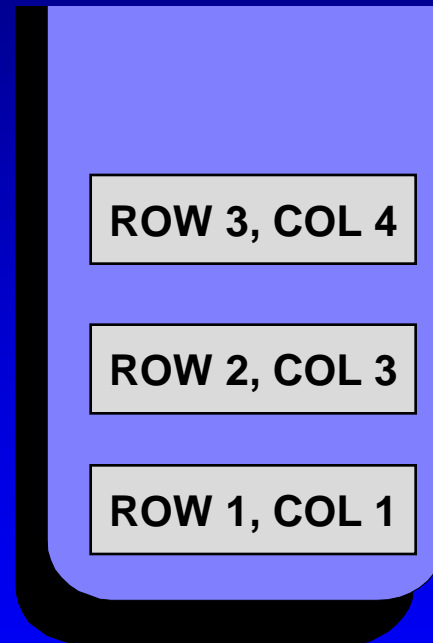
...and we shift to column 4.  
There's still a conflict in column 4, so we try to shift rightward again...



filled

# How the program works

...but there's  
nowhere else to  
go.

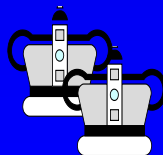
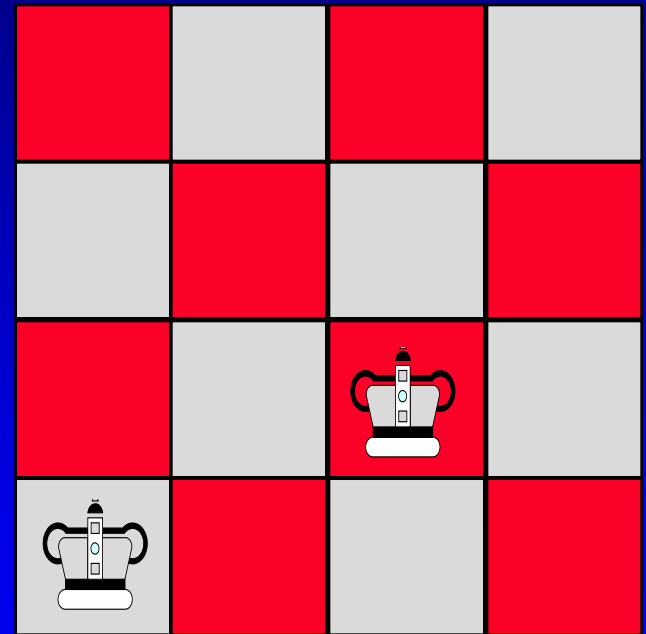
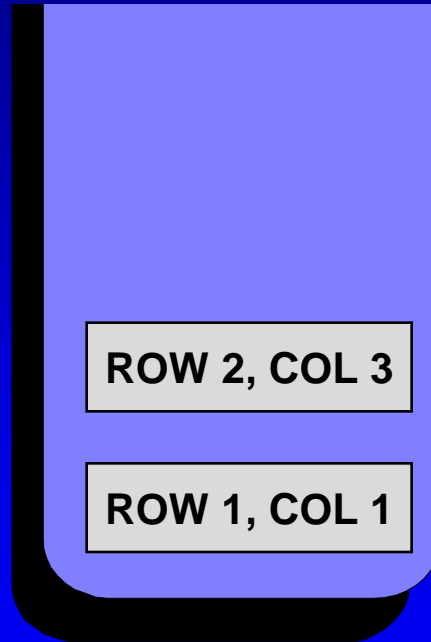


filled

# How the program works

When we run out of room in a row:

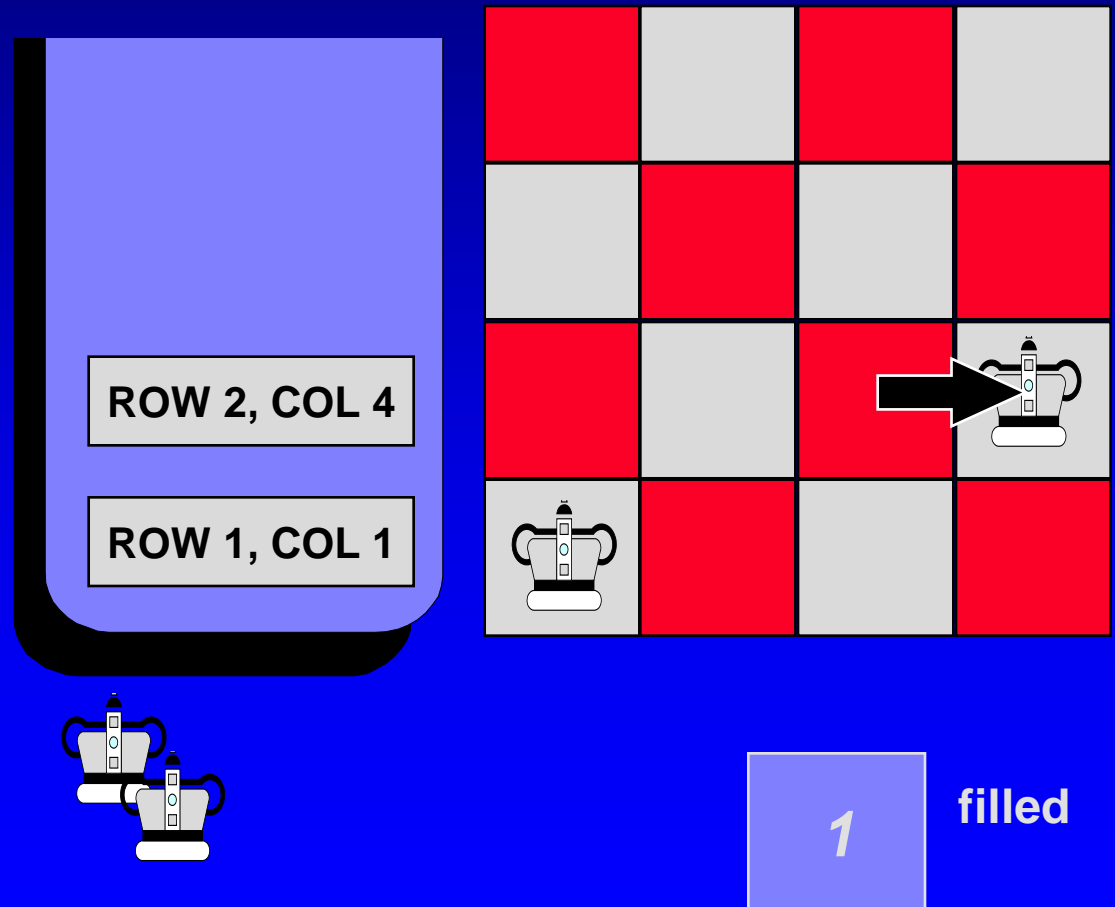
- ❑ pop the stack,
- ❑ reduce filled by 1
- ❑ and continue working on the previous row.



1 filled

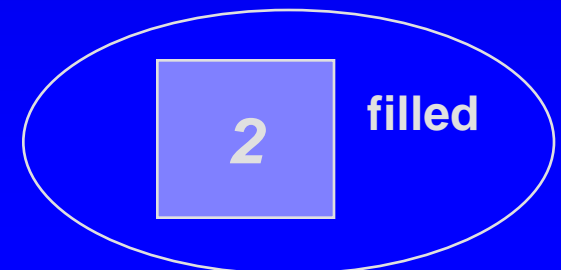
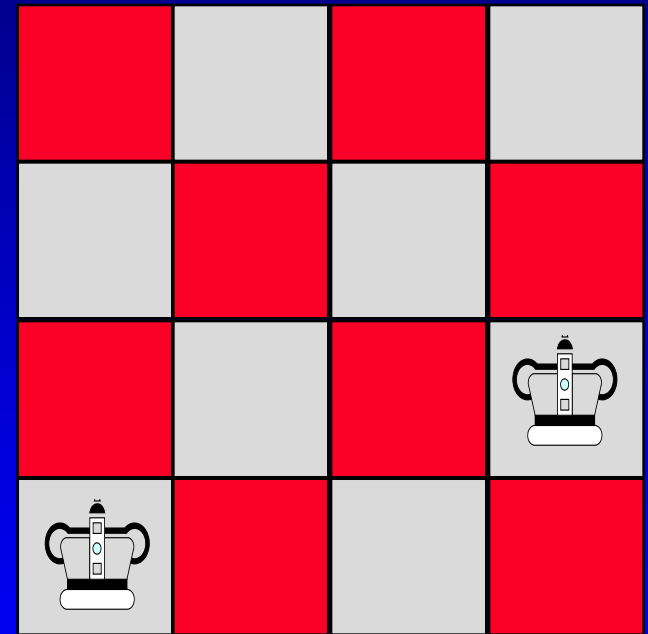
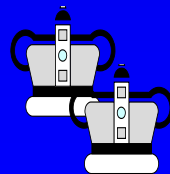
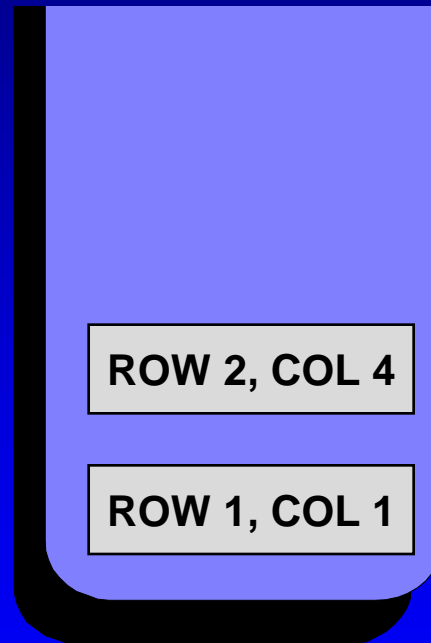
# How the program works

Now we  
continue  
working on row  
2, shifting the  
queen to the  
right.



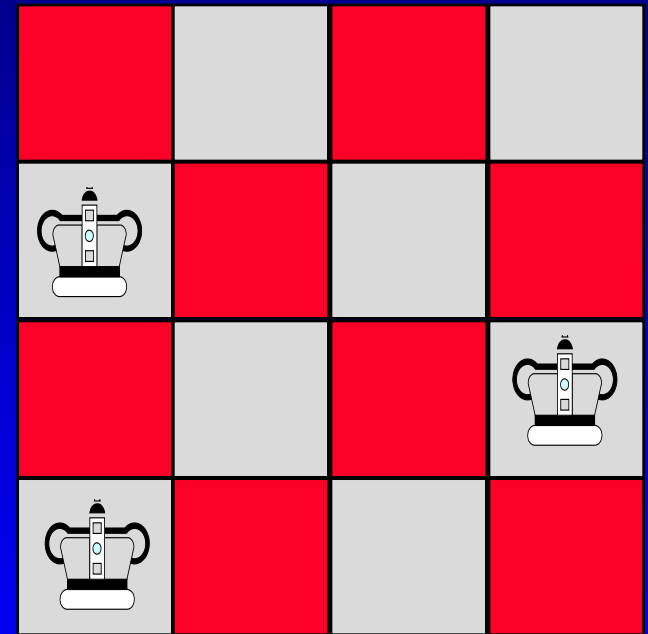
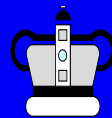
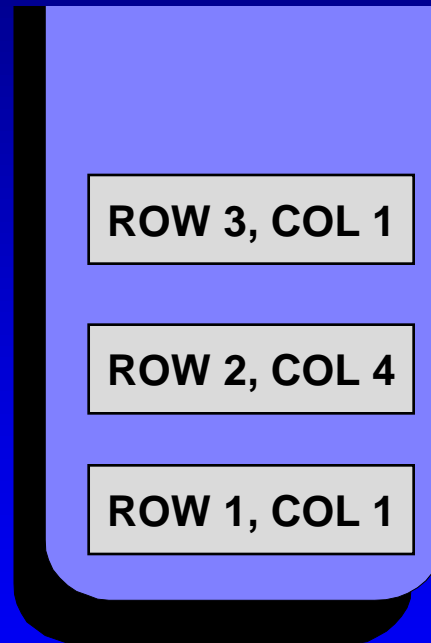
# How the program works

This position has no conflicts, so we can increase filled by 1, and move to row 3.



# How the program works

In row 3, we start again at the first column.



filled



# Pseudocode for N-Queens

- ⌚ Initialize a stack where we can keep track of our decisions.
- Place the first queen, pushing its position onto the stack and setting filled to 0.
- ↺ repeat these steps
  - ▣ if there are no conflicts with the queens...
  - ▣ else if there is a conflict and there is room to shift the current queen rightward...
  - ▣ else if there is a conflict and there is no room to shift the current queen rightward...

# Pseudocode for N-Queens

Ž repeat these steps

□ if there are no conflicts with the queens...

Increase filled by 1. If filled is now N, then the algorithm is done. Otherwise, move to the next row and place a queen in the first column.

# Pseudocode for N-Queens

Ž repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...

Move the current queen rightward,  
adjusting the record on top of the stack  
to indicate the new position.

# Pseudocode for N-Queens

Ž repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...
- else if there is a conflict and there is no room to shift the current queen rightward...

Backtrack!

Keep popping the stack, and reducing filled by 1, until you reach a row where the queen can be shifted rightward. Shift this queen right.

# Pseudocode for N-Queens

Ž repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...
- else if there is a conflict and there is no room to shift the current queen rightward...

Backtrack!

Keep popping the stack, and reducing filled by 1, **until you reach a row where the queen can be shifted rightward.** Shift this queen right.

# Watching the program work

---

You can double  
click the left mouse  
button here to run  
the demonstration  
program a second  
time:





# Summary

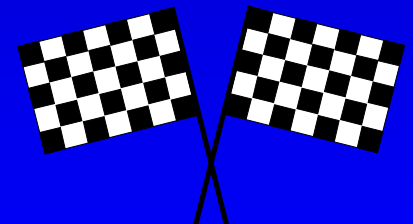
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- ❑ Stacks have many applications.
- ❑ The application which we have shown is called **backtracking**.
- ❑ The key to backtracking: Each choice is recorded in a stack.
- ❑ When you run out of choices for the current decision, you pop the stack, and continue trying different choices for the previous decision.

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THE END