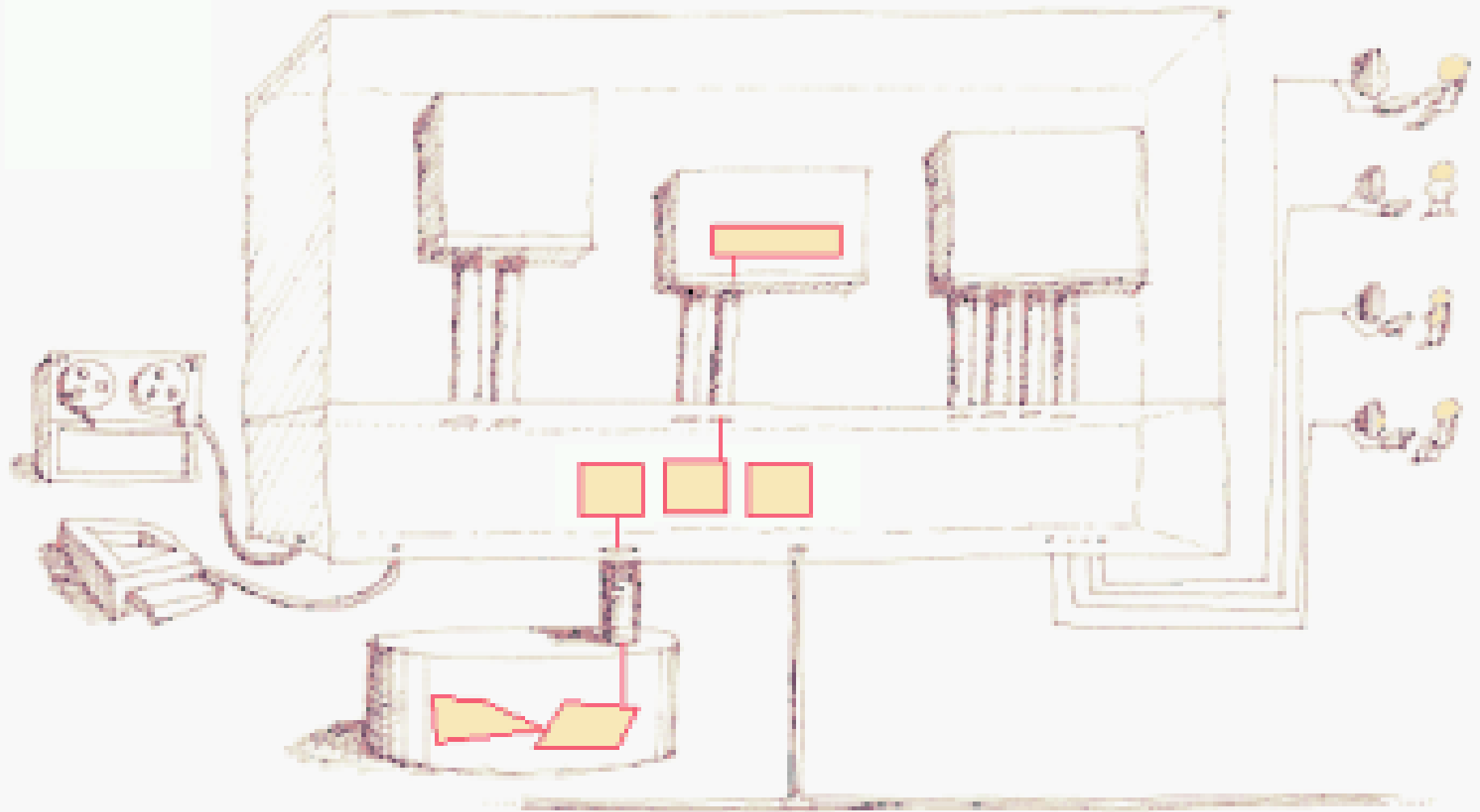


Lecture 2: Users, Files, On-Line Documents

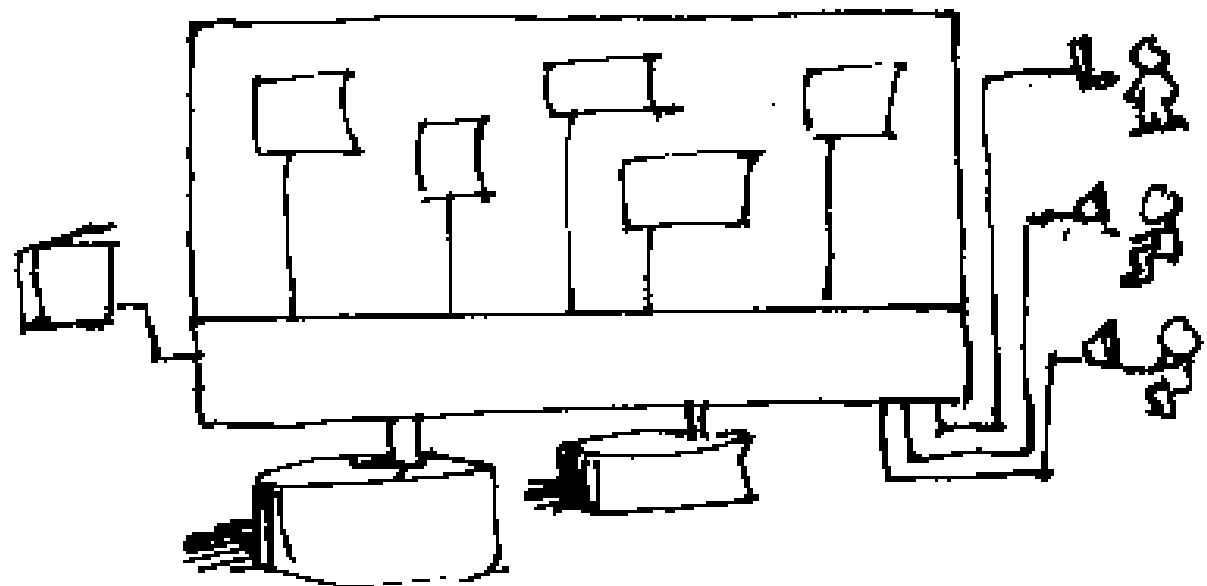


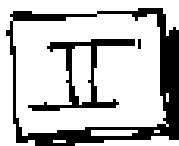
Class 2: Files, Users, On-line Documentation

Summary: Today we shall write a version of the who utility. In the process, we shall learn about:

☆ files, users, time, buffering, AND the fact that Unix is self-documenting. ☆

Recall:





How Does who Work?

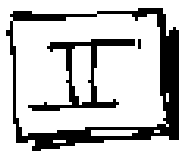
A. What does it do?

purpose:

output:

B. How does it do it?

*C. How can I learn about
the details?*



How Does who Work?

A. What does it do?

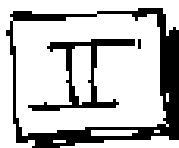
purpose: list users currently
logged on

output: logname, terminal,
time, from where

- * By running the command we see what it does.
- * By consulting the on-line manual..

`man who`

we learn even more.



How Does who Work?

B. How does who do it?

To learn more about Unix commands:

- * read the manual

```
man who
```

- * search the manual

```
man -k user | more
```

- * read the .h files in /usr/include

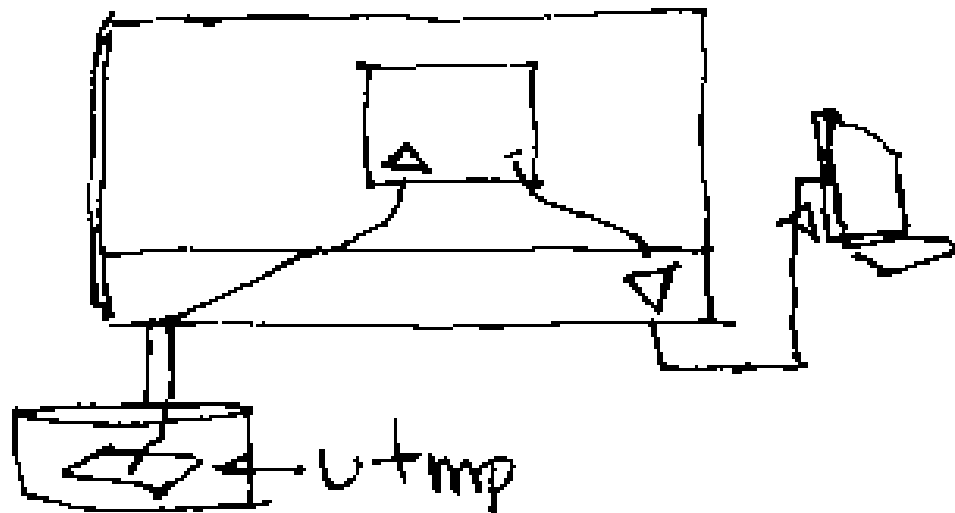
```
more /usr/include/foo.h
```

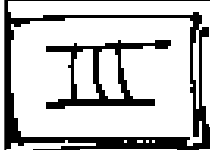
- * follow the "See Also" links

Answer

who works by:

```
open wtmp  
→ read record  
  display info  
  close file  
← eof
```





Q: How Do I Read structs from a File?

If you have used `getc()` and `fgets()`, you know how to read characters and lines, but what about structs of raw data?

What Does the Manual Say?

```
man -k file | grep read
```

```
man 2 read ← topic
```

↑
section

Answer:

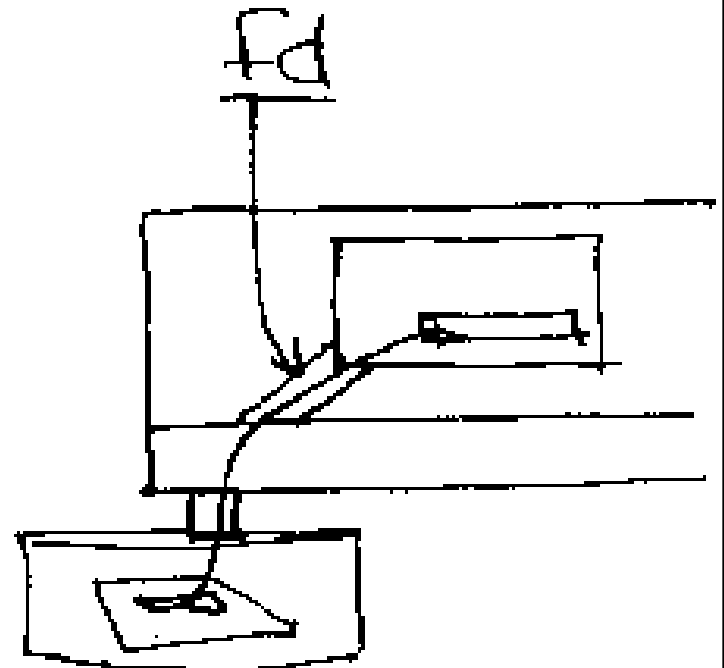
We use `open()`, `read()`, and `close()`.

```
fd = open(name, mode)
```

 ↑ ↙
char * O_RDONLY
 O_WRONLY
 O_RDWR

returns -1 on error
OR an int on success

*creates a
connection to
a file.*



Answer:

We use `open()`, `read()`, and `close()`.

```
fd = open(name, mode)
```

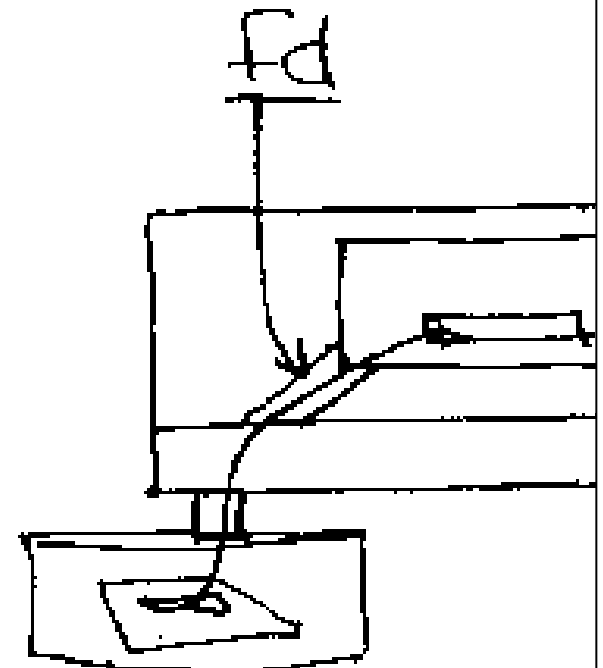
```
n = read(fd, array, numchars)
```

- * transfers numchars from file to the array

- * returns number of chars actually transferred, OR -1 on error

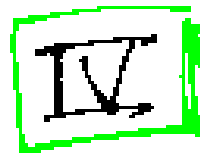
```
close(fd)
```

- * destroys the connection



Now we can write who1.c

*It is getting pretty close,
but it needs some work..*



How Do We Get It to Look Good?

A. Suppress blank entries

B. Format time



Let's Check the Manual



(And the header files)

A. Suppress blank entries

check `ut_type` (see `utmp.h`)

B. Format time

`ctime()` converts to string
Unix stores time as seconds
since the beginning of the
Epoch

Result: `who2.c`

Moral: Unix provides complete
documentation on system
structures and programs

V. Project 2: cp (read and write)

*In who, we read from a file.
How do we write to a file?
Let's explore a real example:*

`cp sourcefile copyfile`

(1) What does cp do?

*ANS: creates or truncates a file,
then writes data into it*

(2) How does cp create and write?

*ANS: search the manual for
the answers*

Creating/Truncating A File

```
fd = creat(name, 0644)
```

First: creates or truncates a file

Then: opens the file for writing

If create, set mode to, 0644

returns -1 on error,

otherwise a file descriptor

Writing to A File

```
n = write(fd, buffer, num)
```

** sends 'num' chars to the file*

** returns actual number sent*

OR returns -1 on error

VI. Does Buffer Size Matter?

Ex: Filesize = 2500 bytes
if buffer = 100 bytes,
=> 25 read() and 25 write() calls

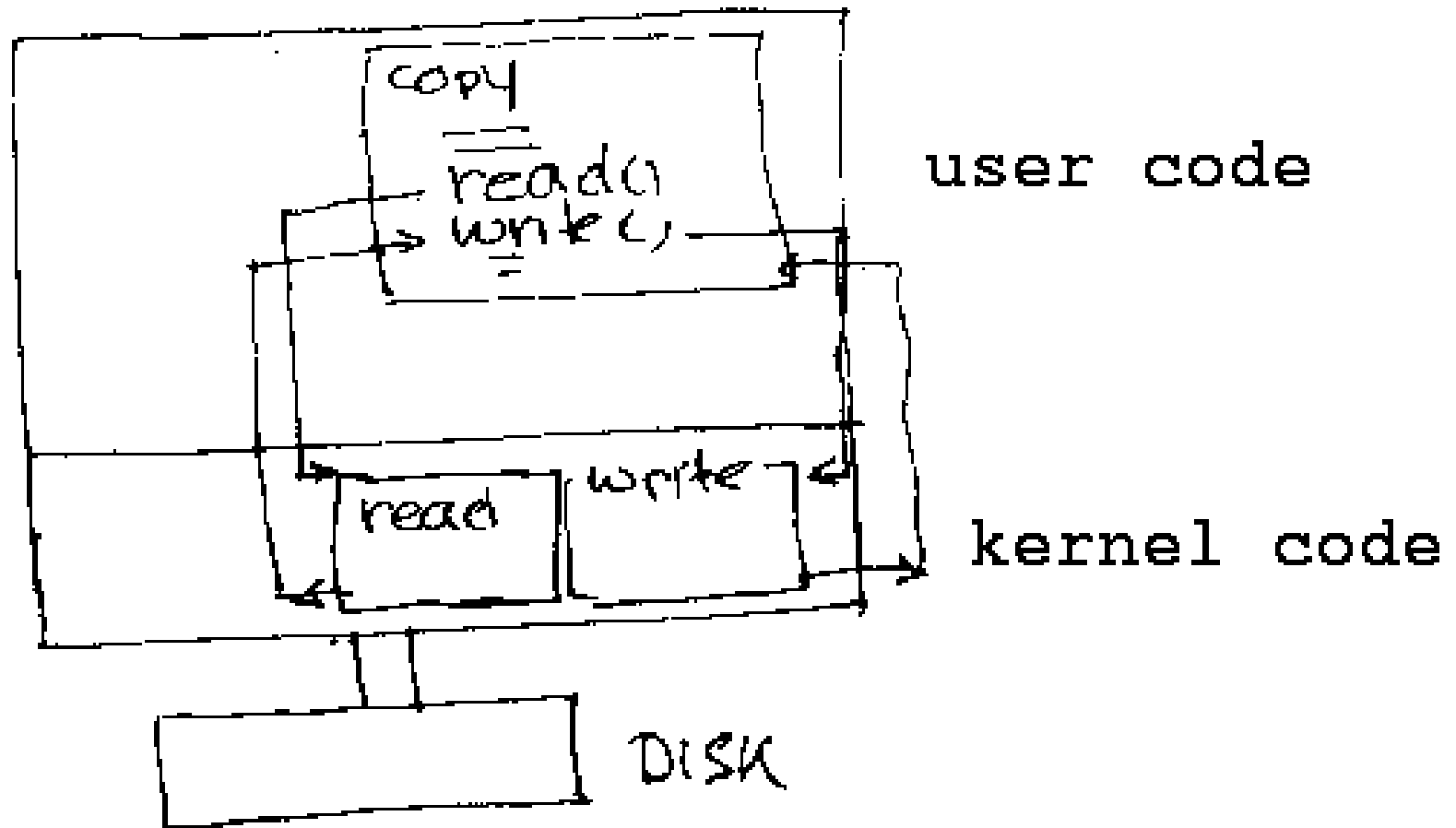
if buffer = 1000 bytes
=> 3 read() and 3 write() calls

Important Idea:

A system call is resource 'expensive' (i.e. takes time). It runs various kernel functions, and it also requires a shift from USER MODE to KERNEL MODE and back.

**** Thus, try to minimize system calls ****

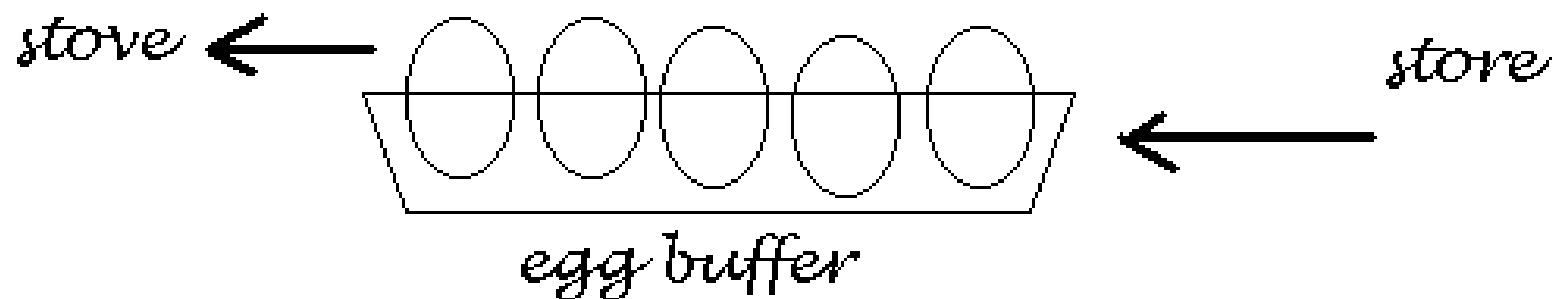
Control flow in copying a file:



VII. Does This Mean `who2.c` Is Inefficient?

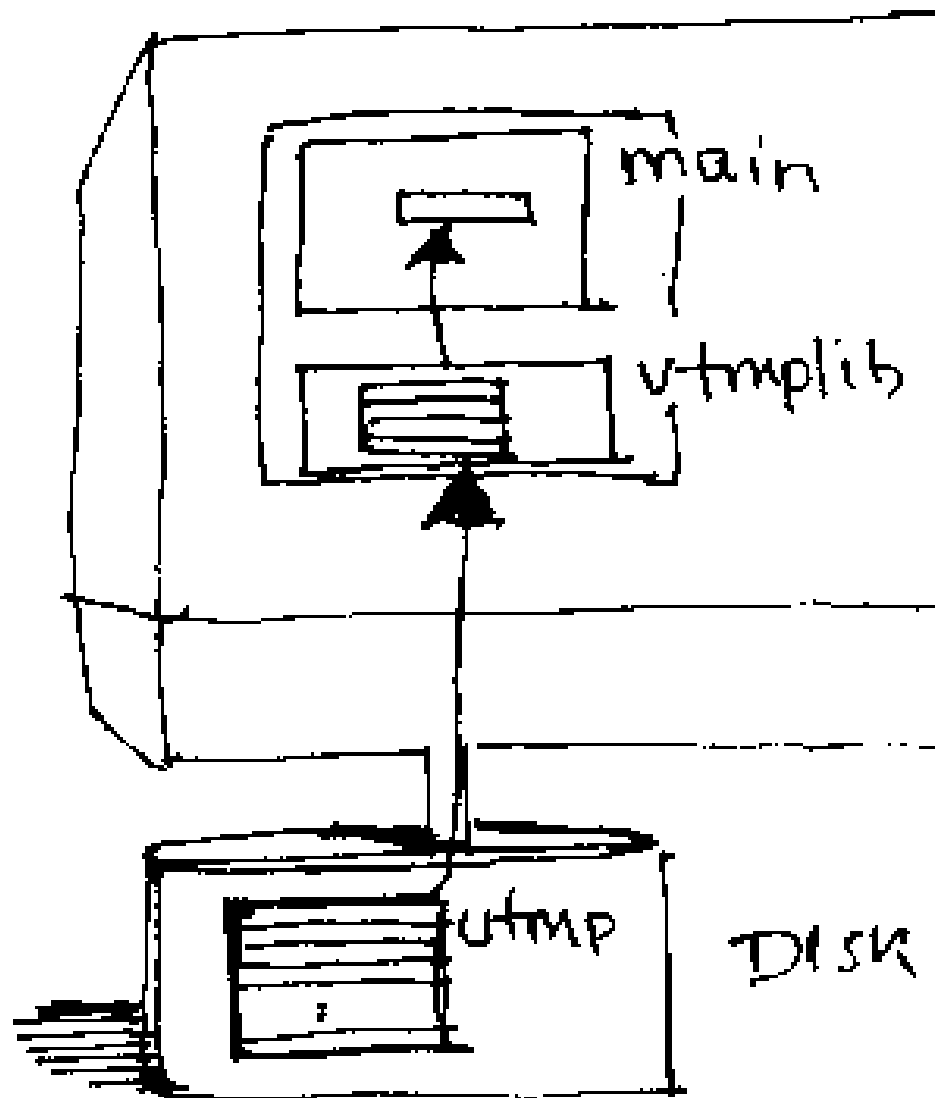
Yes! Making one system call for each line of output makes as much sense as buying pizza by the slice or eggs one at a time.

Better Idea: Read in a bunch of records at a time and then, as with eggs in a carton, take them one by one.



Adding Buffering to who2.c

Picture:



main calls
functions in
utmp.c

utmp.c keeps
a utmp buffer
with space for 16
utmp structs.

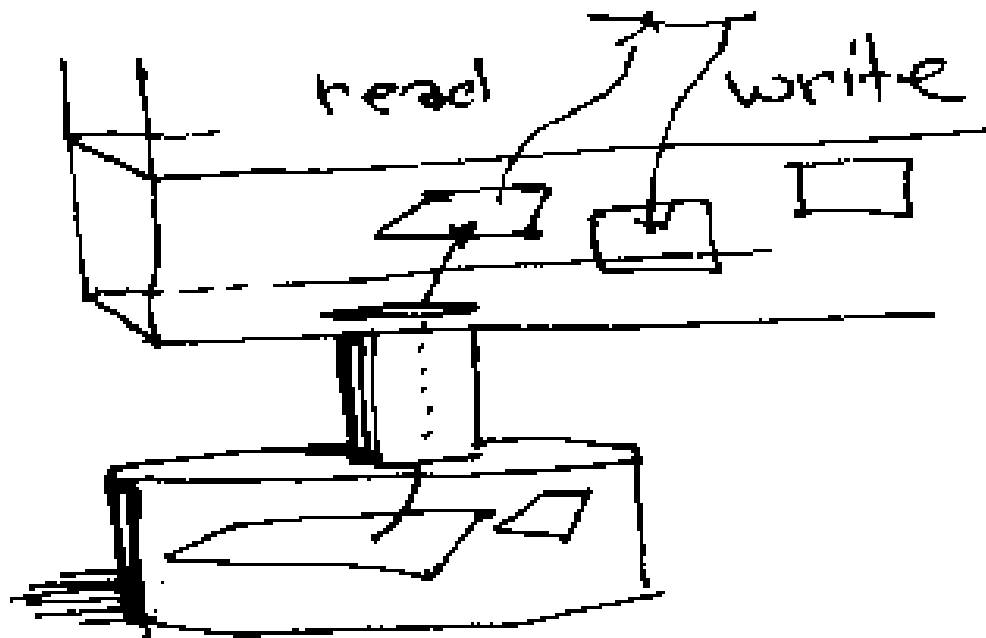
utmp.c reads
only when the
buffer empties.

Result:
who3.c

VIII. If Buffering Is So Smart, Why Doesn't the Kernel Do It?

It does!

The kernel keeps copies of disk blocks in memory. It writes those blocks to disk now and then. The `read()` call actually copies data from kernel buffers, not from the disk.



Consequences of Buffering

- 1) Faster "disk" I/O
- 2) Optimal disk reads and writes
- 3) Need to sync disks before shutdown

IX. Final Example: tail

- 1) What does tail do?
 - 2) How does tail work?
- ...

IX. Final Example: tail


1) What does tail do?

2) How does tail work?

Q: How can a program jump to the end (or to any arbitrary location) of a file?

A: `lseek(fd, offset, base)`

# of chars


0: start,
1: current
2: end

The next read or write occurs there.