

Machine-Level Programming IV: Data

15-213: Introduction to Computer Systems
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Today

■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

■ Structures

- Allocation
- Access
- Alignment

■ Floating Point

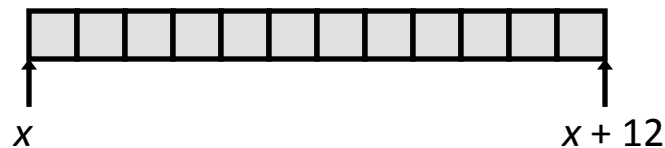
Array Allocation

■ Basic Principle

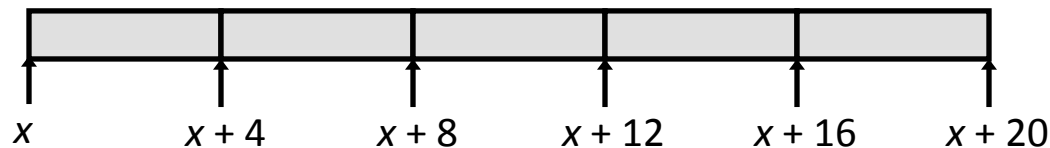
T $A[L]$;

- Array of data type T and length L
- Contiguously allocated region of $L * \text{sizeof}(T)$ bytes in memory

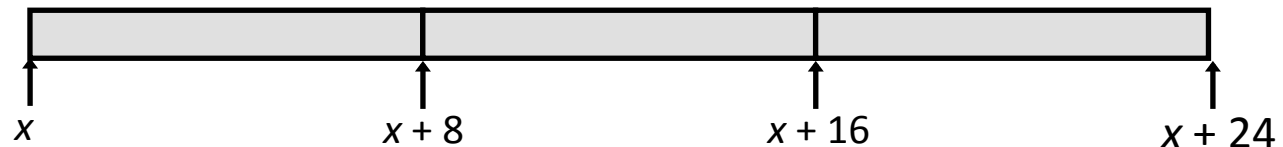
`char string[12];`



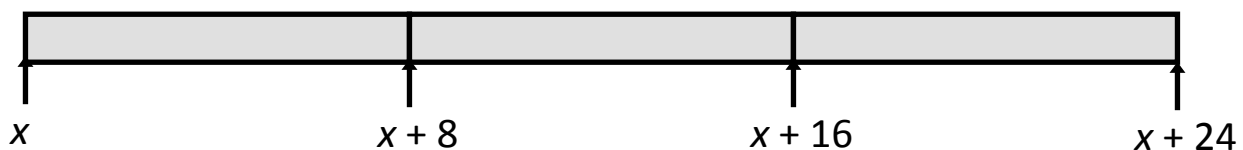
`int val[5];`



`double a[3];`



`char *p[3];`

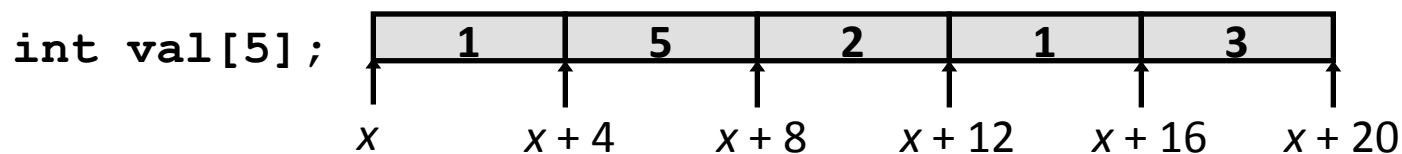


Array Access

■ Basic Principle

T **A**[L];

- Array of data type T and length L
- Identifier **A** can be used as a pointer to array element 0: Type T^*

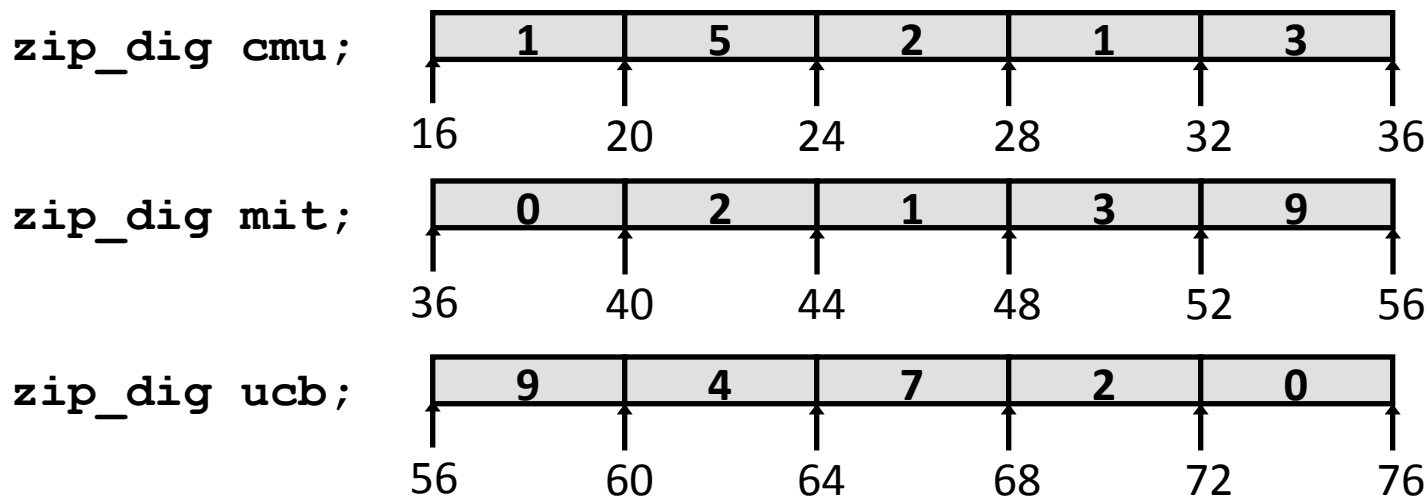


■ Reference	Type	Value
<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	x
<code>val+1</code>	<code>int *</code>	$x + 4$
<code>&val[2]</code>	<code>int *</code>	$x + 8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val + i</code>	<code>int *</code>	$x + 4\ i$

Array Example

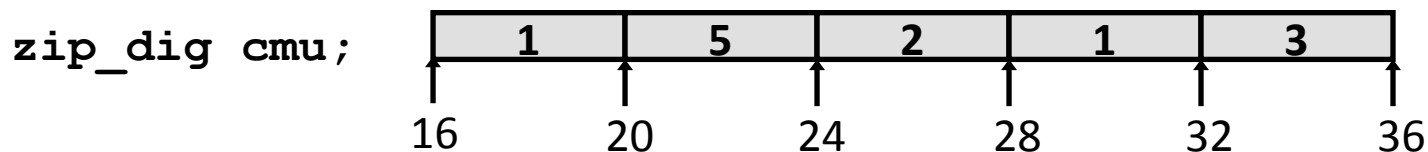
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “`zip_dig cmu`” equivalent to “`int cmu[5]`”
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

IA32

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register `%rdi` contains starting address of array
- Register `%rsi` contains array index
- Desired digit at $4 * \%rdi + \%rsi$
- Use memory reference `(%rdi,%rsi,4)`

Array Loop Example

```
void zincr(zip_dig z) {  
    size_t i;  
    for (i = 0; i < ZLEN; i++)  
        z[i]++;  
}
```

```
# %rdi = z  
movl    $0, %eax           # i = 0  
jmp     .L3                # goto middle  
.L4:                        # loop:  
    addl    $1, (%rdi,%rax,4) # z[i]++  
    addq    $1, %rax        # i++  
.L3:                        # middle  
    cmpq    $4, %rax        # i:4  
    jbe     .L4             # if <=, goto loop  
rep; ret
```

Multidimensional (Nested) Arrays

■ Declaration

$T \ A[R][C];$

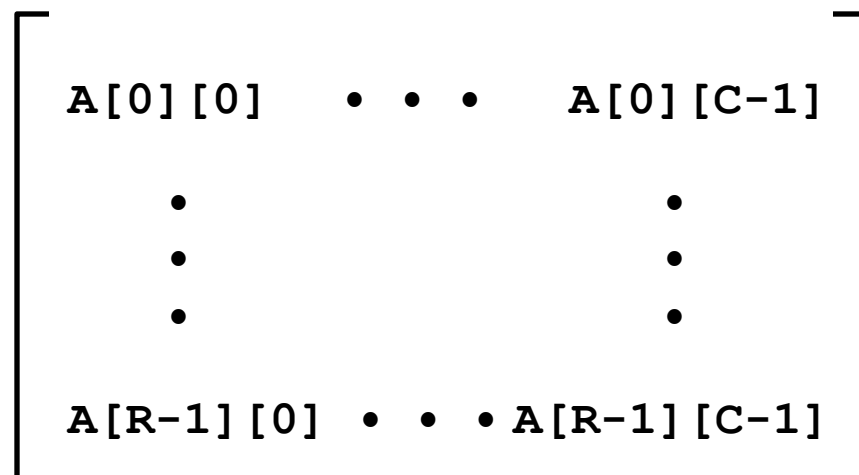
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

■ Array Size

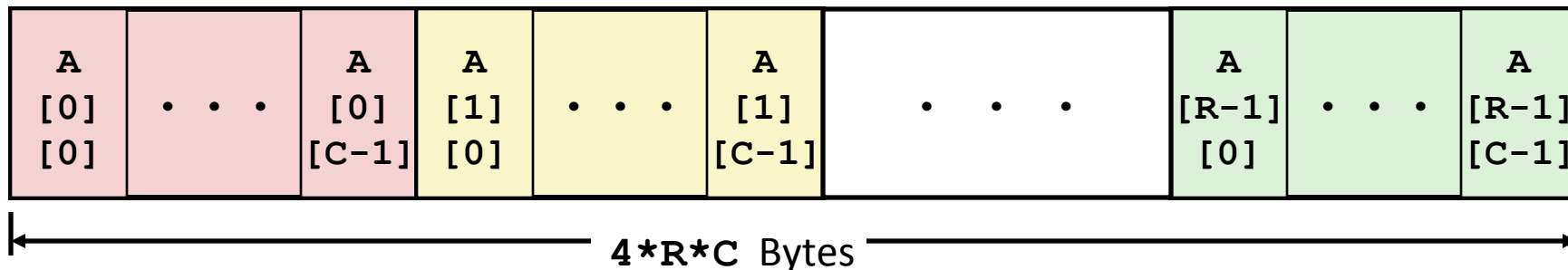
- $R * C * K$ bytes

■ Arrangement

- Row-Major Ordering

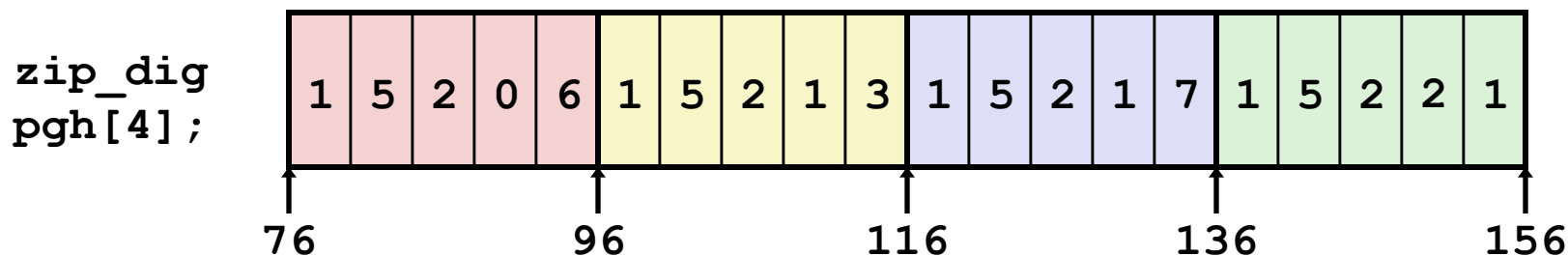


`int A[R][C];`



Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```



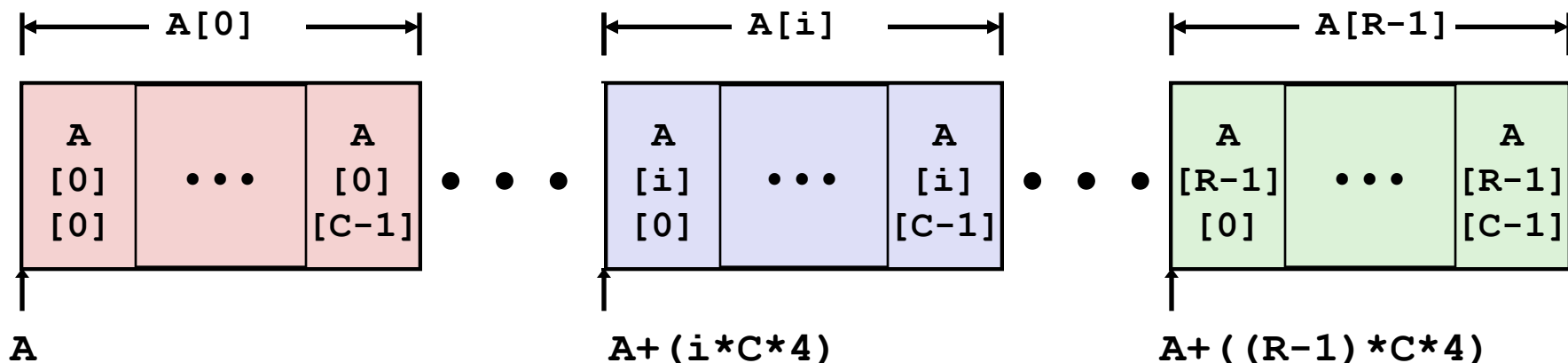
- **“zip_dig pgh[4]” equivalent to “int pgh[4][5]”**
 - Variable **pgh**: array of 4 elements, allocated contiguously
 - Each element is an array of 5 **int**’s, allocated contiguously
- **“Row-Major” ordering of all elements in memory**

Nested Array Row Access

■ Row Vectors

- $\mathbf{A}[i]$ is array of C elements
- Each element of type T requires K bytes
- Starting address $\mathbf{A} + i * (C * K)$

```
int A[R][C];
```

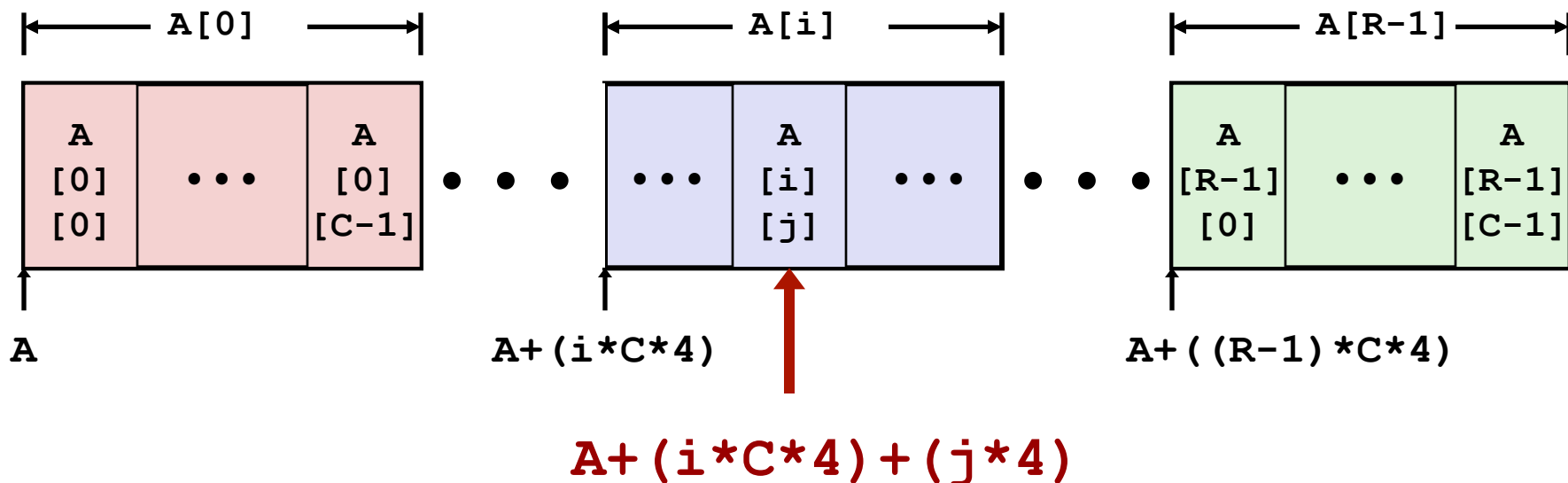


Nested Array Element Access

■ Array Elements

- $A[i][j]$ is element of type T , which requires K bytes
- Address $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```

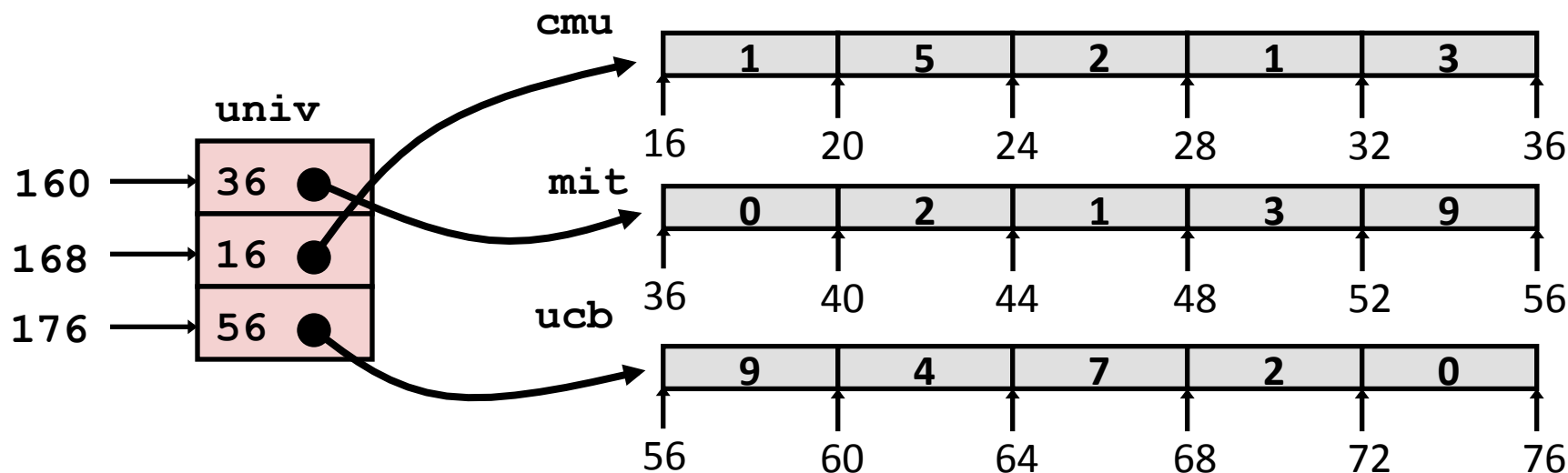


Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig mit = { 0, 2, 1, 3, 9 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3  
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of `int`'s



Element Access in Multi-Level Array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```

```
salq    $2, %rsi                # 4*digit
addq    univ(,%rdi,8), %rsi      # p = univ[index] + 4*digit
movl    (%rsi), %eax             # return *p
ret
```

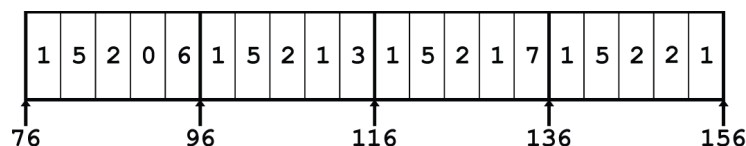
■ Computation

- Element access **Mem[Mem[univ+8*index]+4*digit]**
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

Array Element Accesses

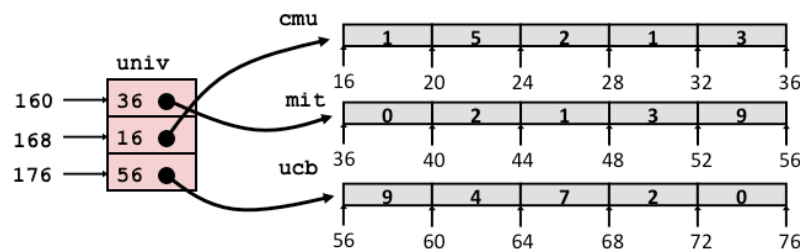
Nested array

```
int get_pgh_digit
(size_t index, size_t digit)
{
    return pgh[index][digit];
}
```



Multi-level array

```
int get_univ_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



Accesses looks similar in C, but address computations very different:

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{digit}]$ $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$

N X N Matrix Code

■ Fixed dimensions

- Know value of N at compile time

```
#define N 16
typedef int fix_matrix[N][N];
/* Get element a[i][j] */
int fix_ele(fix_matrix a,
            size_t i, size_t j)
{
    return a[i][j];
}
```

■ Variable dimensions, explicit indexing

- Traditional way to implement dynamic arrays

```
#define IDX(n, i, j) ((i)*(n)+(j))
/* Get element a[i][j] */
int vec_ele(size_t n, int *a,
            size_t i, size_t j)
{
    return a[IDX(n,i,j)];
}
```

■ Variable dimensions, implicit indexing

- Now supported by gcc

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n],
            size_t i, size_t j) {
    return a[i][j];
}
```

16 X 16 Matrix Access

■ Array Elements

- Address $A + i * (C * K) + j * K$
- $C = 16, K = 4$

```
/* Get element a[i][j] */  
int fix_ele(fix_matrix a, size_t i, size_t j) {  
    return a[i][j];  
}
```

```
# a in %rdi, i in %rsi, j in %rdx  
salq    $6, %rsi          # 64*i  
addq    %rsi, %rdi         # a + 64*i  
movl    (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]  
ret
```


n X n Matrix Access

■ Array Elements

- Address $A + i * (C * K) + j * K$
- $C = n, K = 4$
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
    return a[i][j];
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq    %rdx, %rdi          # n*i
leaq     (%rsi,%rdi,4), %rax  # a + 4*n*i
movl     (%rax,%rcx,4), %eax  # a + 4*n*i + 4*j
ret
```

Today

■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

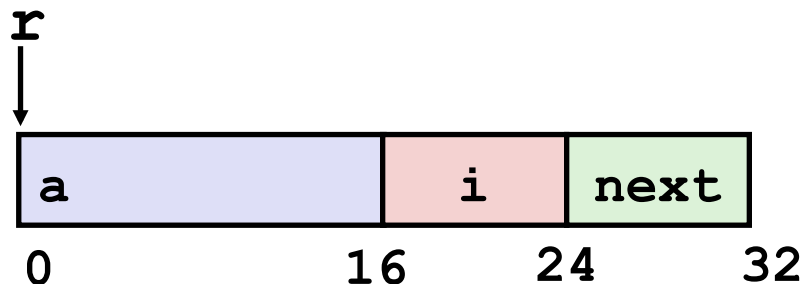
■ Structures

- Allocation
- Access
- Alignment

■ Floating Point

Structure Representation

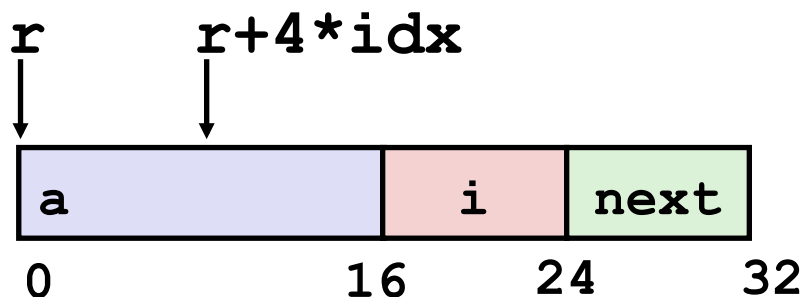
```
struct rec {  
    int a[4];  
    size_t i;  
    struct rec *next;  
};
```



- **Structure represented as block of memory**
 - Big enough to hold all of the fields
- **Fields ordered according to declaration**
 - Even if another ordering could yield a more compact representation
- **Compiler determines overall size + positions of fields**
 - Machine-level program has no understanding of the structures in the source code

Generating Pointer to Structure Member

```
struct rec {
    int a[4];
    size_t i;
    struct rec *next;
};
```



■ Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as $r + 4 \cdot idx$

```
int *get_ap
(struct rec *r, size_t idx)
{
    return &r->a[idx];
}
```

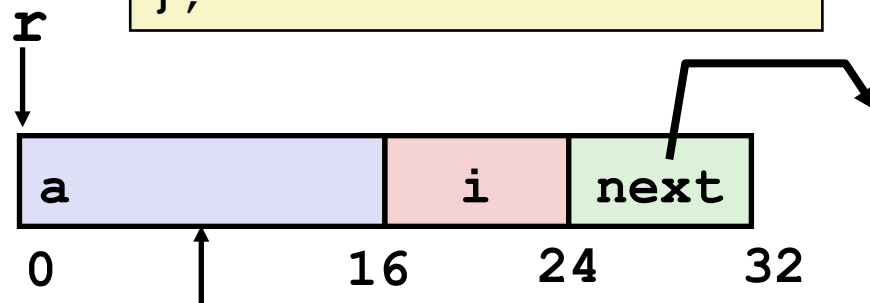
```
# r in %rdi, idx in %rsi
leaq (%rdi,%rsi,4), %rax
ret
```

Following Linked List

■ C Code

```
void set_val
(struct rec *r, int val)
{
    while (r) {
        int i = r->i;
        r->a[i] = val;
        r = r->next;
    }
}
```

```
struct rec {
    int a[3];
    int i;
    struct rec *next;
};
```



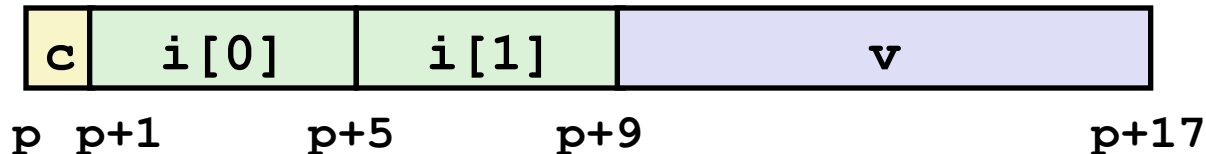
Element `i`

Register	Value
<code>%rdi</code>	<code>r</code>
<code>%rsi</code>	<code>val</code>

```
.L11:                                # loop:
    movslq    16(%rdi), %rax          # i = M[r+16]
    movl      %esi, (%rdi,%rax,4)     # M[r+4*i] = val
    movq      24(%rdi), %rdi         # r = M[r+24]
    testq     %rdi, %rdi             # Test r
    jne       .L11                  # if !=0 goto loop
```

Structures & Alignment

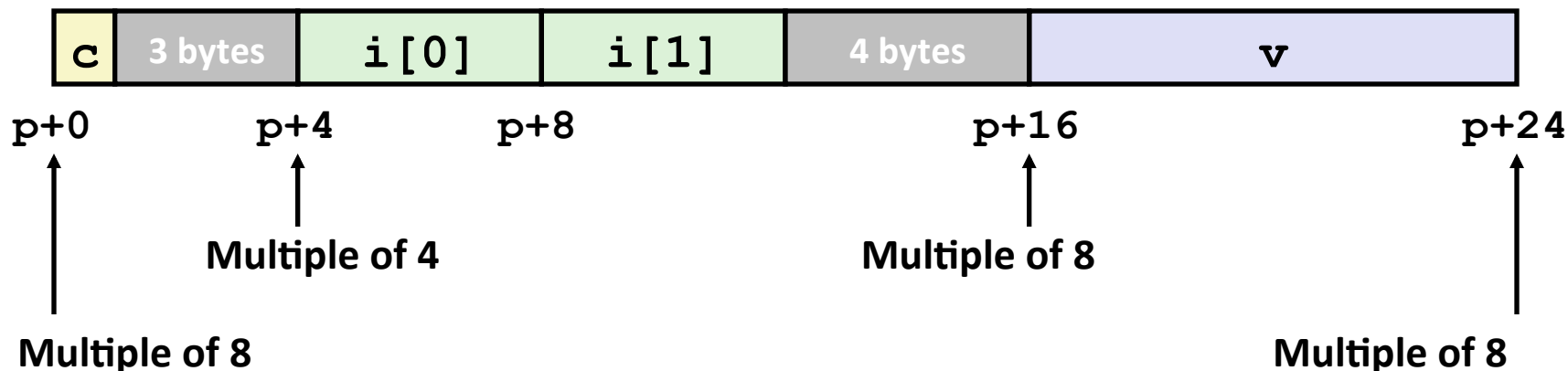
■ Unaligned Data



```
struct S1 {
    char c;
    int i[2];
    double v;
} *p;
```

■ Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



Alignment Principles

■ Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K
- Required on some machines; advised on x86-64

■ Motivation for Aligning Data

- Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
 - Inefficient to load or store datum that spans quad word boundaries
 - Virtual memory trickier when datum spans 2 pages

■ Compiler

- Inserts gaps in structure to ensure correct alignment of fields

Specific Cases of Alignment (x86-64)

- **1 byte: `char`, ...**
 - no restrictions on address
- **2 bytes: `short`, ...**
 - lowest 1 bit of address must be 0_2
- **4 bytes: `int`, `float`, ...**
 - lowest 2 bits of address must be 00_2
- **8 bytes: `double`, `long`, `char *`, ...**
 - lowest 3 bits of address must be 000_2
- **16 bytes: `long double` (GCC on Linux)**
 - lowest 4 bits of address must be 0000_2

Satisfying Alignment with Structures

■ Within structure:

- Must satisfy each element's alignment requirement

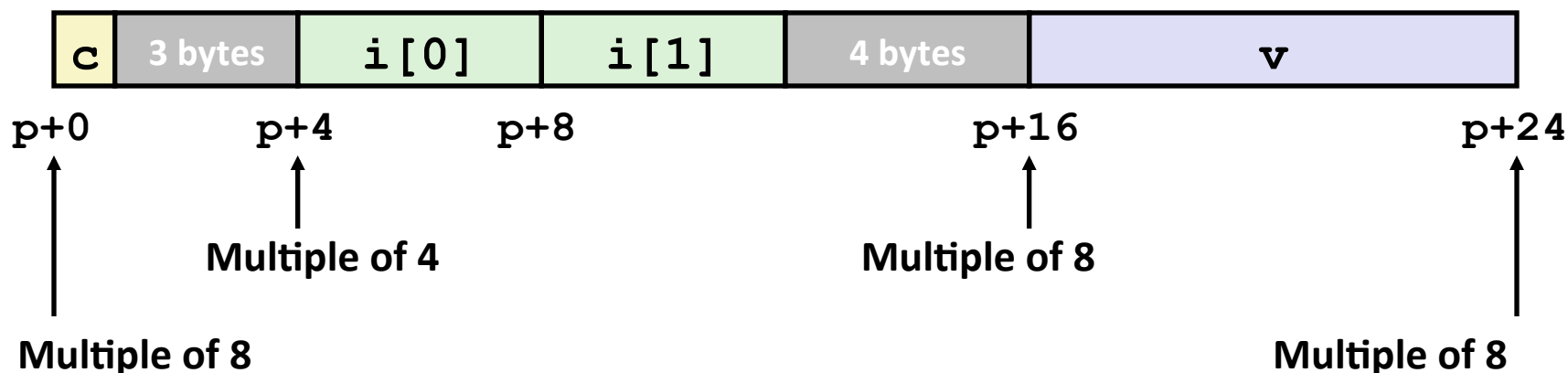
■ Overall structure placement

- Each structure has alignment requirement K
 - K = Largest alignment of any element
- Initial address & structure length must be multiples of K

```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *p;
```

■ Example:

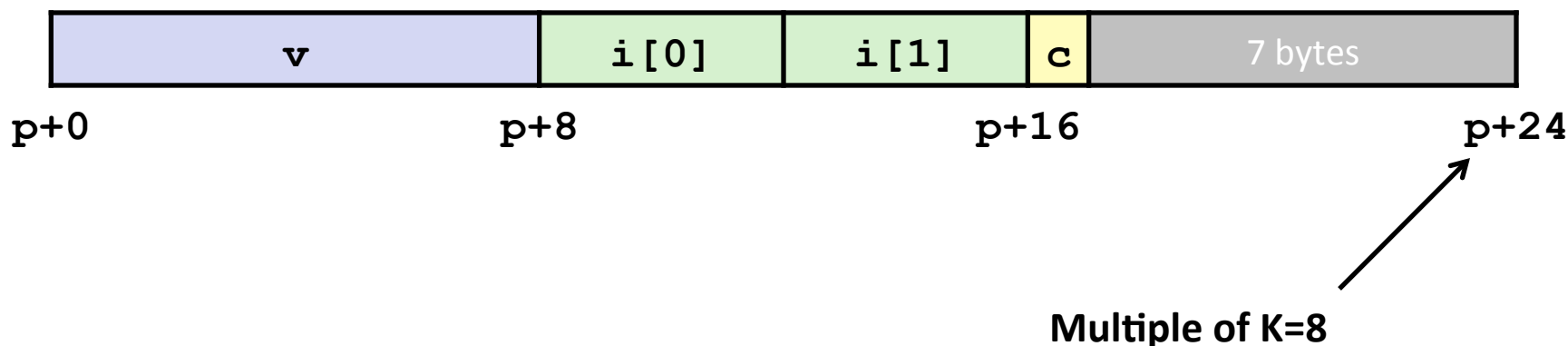
- $K = 8$, due to `double` element



Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

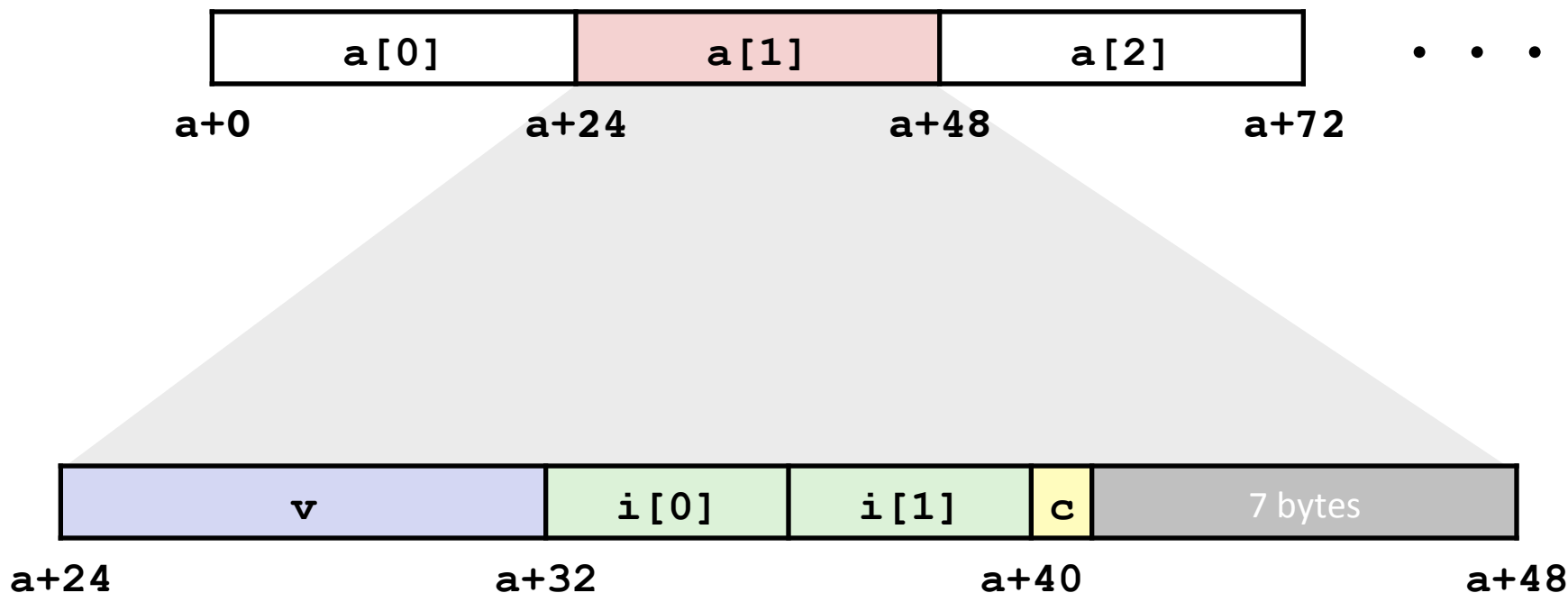
```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} *p;
```



Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

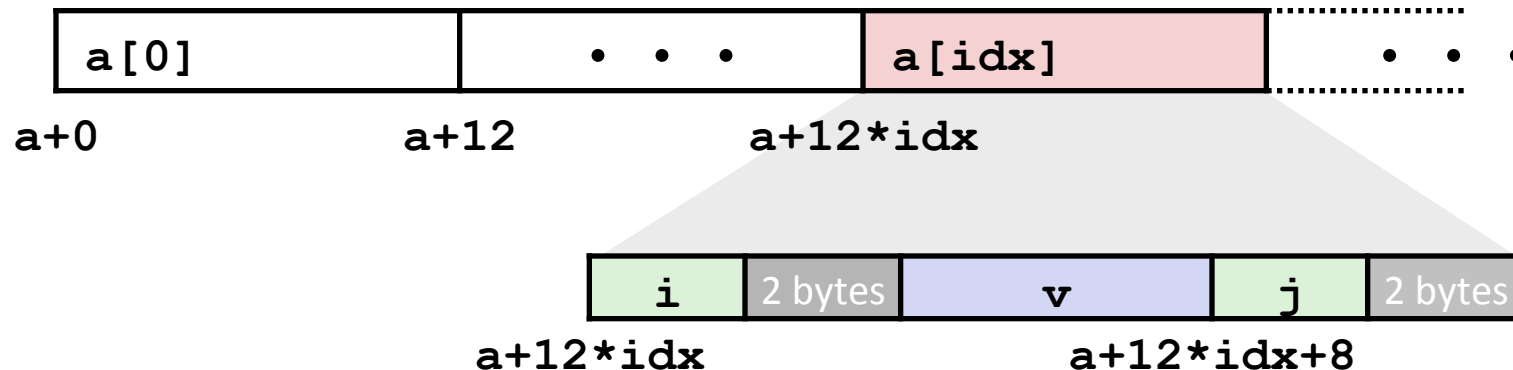
```
struct S2 {  
    double v;  
    int i[2];  
    char c;  
} a[10];
```



Accessing Array Elements

- **Compute array offset $12 \cdot \text{idx}$**
 - `sizeof(S3)`, including alignment spacers
- **Element `j` is at offset 8 within structure**
- **Assembler gives offset `a+8`**
 - Resolved during linking

```
struct S3 {
    short i;
    float v;
    short j;
} a[10];
```



```
short get_j(int idx)
{
    return a[idx].j;
}
```

```
# %rdi = idx
leaq (%rdi,%rdi,2),%rax # 3*idx
movzwl a+8(,%rax,4),%eax
```

Saving Space

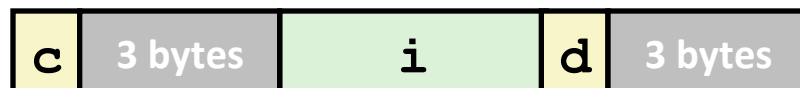
- Put large data types first

```
struct S4 {  
    char c;  
    int i;  
    char d;  
} *p;
```



```
struct S5 {  
    int i;  
    char c;  
    char d;  
} *p;
```

- Effect (K=4)



Today

■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

■ Structures

- Allocation
- Access
- Alignment

■ Floating Point

Background

■ History

- x87 FP
 - Legacy, very ugly
- SSE FP
 - Supported by Shark machines
 - Special case use of vector instructions
- AVX FP
 - Newest version
 - Similar to SSE
 - Documented in book

Programming with SSE3

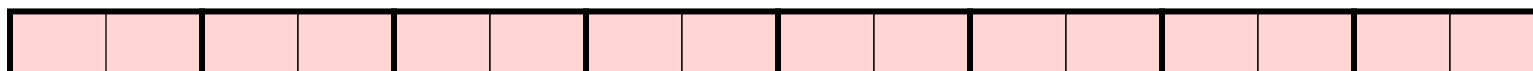
XMM Registers

■ 16 total, each 16 bytes

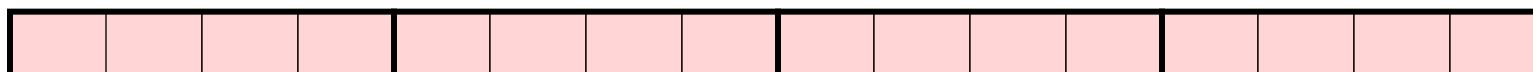
■ 16 single-byte integers



■ 8 16-bit integers



■ 4 32-bit integers



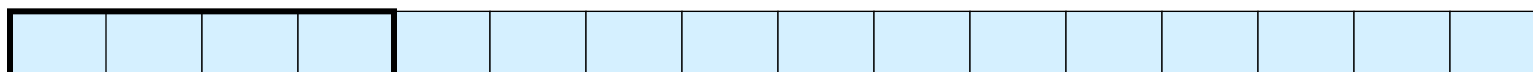
■ 4 single-precision floats



■ 2 double-precision floats



■ 1 single-precision float



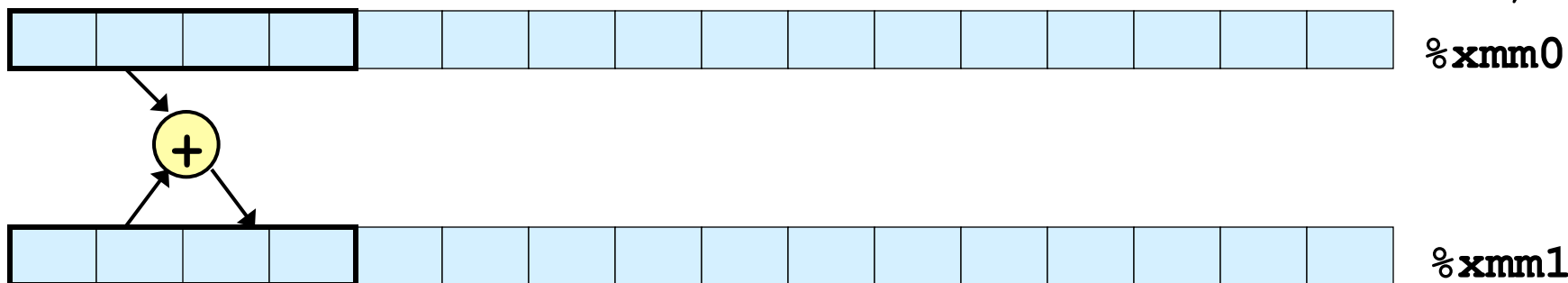
■ 1 double-precision float



Scalar & SIMD Operations

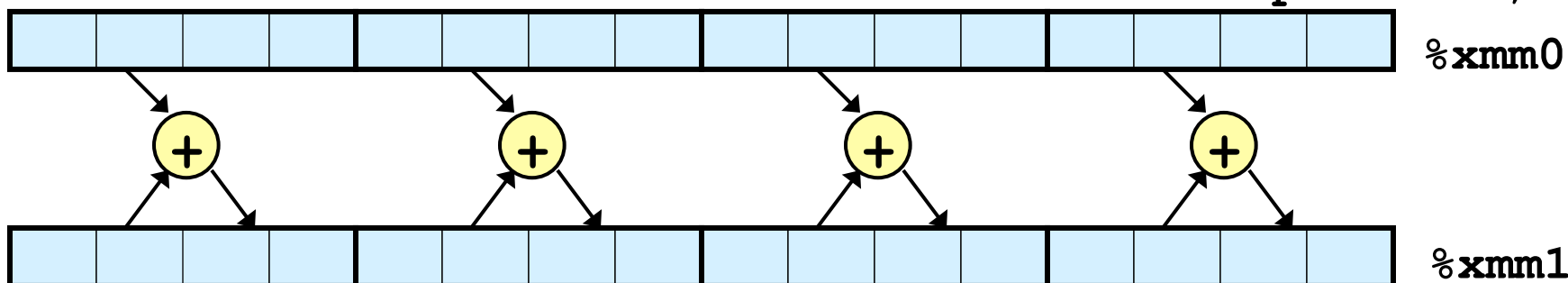
■ Scalar Operations: Single Precision

`addss %xmm0, %xmm1`



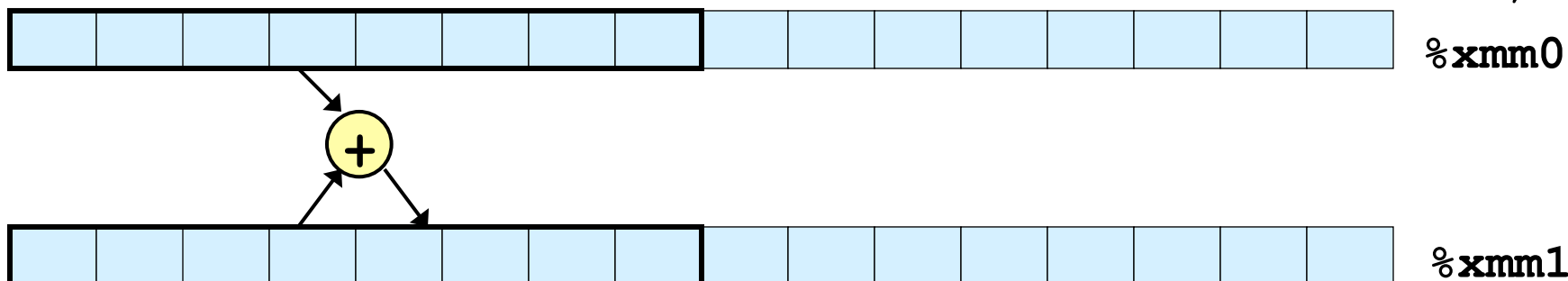
■ SIMD Operations: Single Precision

`addps %xmm0, %xmm1`



■ Scalar Operations: Double Precision

`addsd %xmm0, %xmm1`



FP Basics

- Arguments passed in `%xmm0`, `%xmm1`, ...
- Result returned in `%xmm0`
- All XMM registers caller-saved

```
float fadd(float x, float y)
{
    return x + y;
}
```

```
double dadd(double x, double y)
{
    return x + y;
}
```

```
# x in %xmm0, y in %xmm1
addss    %xmm1, %xmm0
ret
```

```
# x in %xmm0, y in %xmm1
addsd    %xmm1, %xmm0
ret
```

FP Memory Referencing

- Integer (and pointer) arguments passed in regular registers
- FP values passed in XMM registers
- Different mov instructions to move between XMM registers, and between memory and XMM registers

```
double dincr(double *p, double v)
{
    double x = *p;
    *p = x + v;
    return x;
}
```

```
# p in %rdi, v in %xmm0
movapd  %xmm0, %xmm1    # Copy v
movsd   (%rdi), %xmm0    # x = *p
addsd   %xmm0, %xmm1     # t = x + v
movsd   %xmm1, (%rdi)    # *p = t
ret
```

Other Aspects of FP Code

■ *Lots of instructions*

- Different operations, different formats, ...

■ Floating-point comparisons

- Instructions `ucomiss` and `ucomisd`
- Set condition codes CF, ZF, and PF

■ Using constant values

- Set XMM0 register to 0 with instruction `xorpd %xmm0, %xmm0`
- Others loaded from memory

Summary

■ Arrays

- Elements packed into contiguous region of memory
- Use index arithmetic to locate individual elements

■ Structures

- Elements packed into single region of memory
- Access using offsets determined by compiler
- Possible require internal and external padding to ensure alignment

■ Combinations

- Can nest structure and array code arbitrarily

■ Floating Point

- Data held and operated on in XMM registers